

Refutation of Bell's theorem for temporal logic

© Copyright 2019 by Colin James III All rights reserved.

Abstract: We evaluate Bell's theorem for temporal logic. It is *not* tautologous. Hence schematics of a protocol for a violation of Bell's inequalities for temporal order are similarly moot. These conjectures form a *non* tautologous fragment of the universal logic $\forall\mathcal{L}4$.

We assume the method and apparatus of Meth8/ $\forall\mathcal{L}4$ with Tautology as the designated proof value, \mathbf{F} as contradiction, \mathbf{N} as truthity (non-contingency), and \mathbf{C} as falsity (contingency). The 16-valued truth table is row-major and horizontal, or repeating fragments of 128-tables, sometimes with table counts, for more variables. (See ersatz-systems.com.)

LET \sim Not, \neg ; $+$ Or, \vee, \cup, \sqcup ; $-$ Not Or; $\&$ And, $\wedge, \cap, \square, \cdot, \otimes$; \backslash Not And;
 $>$ Imply, greater than, $\rightarrow, \Rightarrow, \mapsto, \succ, \supset, \Rightarrow$; $<$ Not Imply, less than, $\in, <, \subset, \neq, \neq, \leftarrow, \lesssim$;
 $=$ Equivalent, $\equiv, :=, \Leftrightarrow, \leftrightarrow, \triangleq, \approx, \cong$; $@$ Not Equivalent, \neq, \oplus ;
 $\%$ possibility, for one or some, $\exists, \diamond, \mathbf{M}$; $\#$ necessity, for every or all, $\forall, \square, \mathbf{L}$;
 $(z=z)$ \mathbf{T} as tautology, \mathbf{T} , ordinal 3; $(z@z)$ \mathbf{F} as contradiction, $\emptyset, \text{Null}, \perp$, zero;
 $(\%z\>\#z)$ \mathbf{N} as non-contingency, Δ , ordinal 1; $(\%z\<\#z)$ \mathbf{C} as contingency, ∇ , ordinal 2;
 $\sim(y < x)$ ($x \leq y$), ($x \subseteq y$), ($x \sqsubseteq y$); $(A=B)$ ($A\sim B$).
 Note for clarity, we usually distribute quantifiers onto each designated variable.

From: Zych, M.; Costa, F.; Pikovski, I.; Brukner, Č. (2019). Bell's theorem for temporal order. nature.com/articles/s41467-019-11579-x, arxiv.org/pdf/1708.00248.pdf m.zych@uq.edu.au

B. Bell's theorem for temporal order The scenario for which the theorem is formulated involves a bipartite system ...

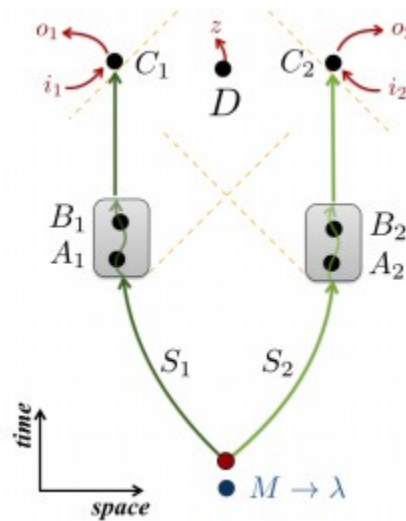


FIG. 2: Bell's theorem for temporal order. A bipartite system, made of subsystems S_1 and S_2 , is sent to two groups of agents. Operations on S_1 (S_2) are performed at events A_1, B_1 (A_2, B_2). At event C_1 (C_2), a measurement with setting i_1 (i_2) and outcome o_1 (o_2) is performed. Events A_1, B_1 are space-like separated from A_2, B_2 , and C_1 is space-like to C_2 ; light cones are marked by dashed yellow lines. ... The system M is measured at event D , producing an output bit z . If the initial state of the systems S_1, S_2, M is separable, and λ is a classical variable ... , the resulting bipartite statistics of the outcomes o_1, o_2 cannot violate any Bell inequality, even if conditioned on z . (1.0)

LET p, q, r, s, t, u, v, w: A, B, C, S, D, i, o, z.

$$\begin{aligned}
 & ((s \& (\%s \> \#s)) \> (((p \& (\%s \> \#s)) \& (q \& (\%s \> \#s))) \> ((r \& (\%s \> \#s)) \> (((u \& (\%s \> \#s)) \> (v \& (\%s \> \#s))) \> \\
 & (t \> z)))))) = \\
 & ((s \& \sim (\%s \> \#s)) \> (((p \& \sim (\%s \> \#s)) \& (q \& \sim (\%s \> \#s))) \> ((r \& \sim (\%s \> \#s)) \> (((u \& \sim (\%s \> \#s)) \> (v \& \sim (\% \\
 & s \> \#s))) \> (t \> z)))))) ;
 \end{aligned}$$

$$\begin{aligned}
 & \text{TTTT TTTT TTTT TTTT (1) } \times 4 \\
 & \text{TTTT TTTT TTTT TTTT} \mathbf{F} \text{(1) } \} \\
 & \text{TTTT TTTT TTTT TTTT (2) } \} \\
 & \text{TTTT TTTT TTTT TTTT} \mathbf{F} \text{(1) } \} \times 3 \} \\
 & \text{TTTT TTTT TTTT TTTT (1) } \} \} \\
 & \text{TTTT TTTT TTTT TTTT (2) } \} \\
 & \text{TTTT TTTT TTTT TTTT (1) } \} \times 2 \} \\
 & \text{TTTT TTTT TTTT TTTT} \mathbf{F} \text{(1) } \} \} \\
 & \text{TTTT TTTT TTTT TTTT (64) } \qquad \qquad \qquad (1.2)
 \end{aligned}$$

Eq. 1.2 as rendered is *not* tautologous. Hence schematics of a protocol for a violation of Bell's inequalities for temporal order are similarly moot.