

Refutation of abductive repair in ontology engineering

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Abstract: We evaluate the stated example of complete-debug problem (CDP) in formulas framing the definitions, oracles, and repairs. None is tautologous. This forms a *non* tautologous fragment of the universal logic VL4.

We assume the method and apparatus of Meth8/VL4 with Tautology as the designated proof value, **F** as contradiction, **N** as truthity (non-contingency), and **C** as falsity (contingency). The 16-valued truth table is row-major and horizontal, or repeating fragments of 128-tables, sometimes with table counts, for more variables. (See ersatz-systems.com.)

LET \sim Not, \neg ; + Or, \vee , \cup , \sqcup ; - Not Or; & And, \wedge , \cap , \square , \cdot , \otimes ; \ Not And;
> Imply, greater than, \rightarrow , \Rightarrow , \mapsto , $>$, \supset , \rightarrow ; < Not Imply, less than, \in , $<$, \subset , \neq , \neq , \ll , \leq ;
= Equivalent, \equiv , $:=$, \Leftrightarrow , \leftrightarrow , $\hat{=}$, \approx , \simeq ; @ Not Equivalent, \neq , \oplus ;
% possibility, for one or some, \exists , \diamond , **M**; # necessity, for every or all, \forall , \square , **L**;
(z=z) **T** as tautology, **T**, ordinal 3; (z@z) **F** as contradiction, \emptyset , Null, \perp , zero;
(%z>#z) **N** as non-contingency, Δ , ordinal 1; (%z<#z) **C** as contingency, ∇ , ordinal 2;
 $\sim(y < x)$ ($x \leq y$), ($x \subseteq y$), ($x \sqsubseteq y$); (A=B) (A~B).
Note for clarity, we usually distribute quantifiers onto each designated variable.

From: Lambrix, P. (2019). Completing and debugging ontologies: state of the art and challenges.
arxiv.org/pdf/1908.03171.pdf

Abstract As semantically-enabled applications require high-quality ontologies, developing and maintaining as correct and complete as possible ontologies is an important, although difficult task in ontology engineering. A key step is ontology debugging and completion. In general, there are two steps: detecting defects and repairing defects. In this paper we formalize the repairing step as an abduction problem and situate the state of the art with respect to this framework. ...

2.1 Formalization

2.1.1 Repair ... As an example, consider the CDP [complete-debug problem] in Fig. 1. ... Then R1,R2, R3, R4 and R5 are all repairs of the CDP

Figure 1: Example complete-debug problem

$$\text{T: } \{ax1: p1 \sqsubseteq p2, ax2: p1 \sqsubseteq p3, ax3: p1 \sqsubseteq \neg p4, ax4: p2 \sqsubseteq p4, ax5: p2 \sqsubseteq p5, ax6: p3 \sqsubseteq p5, \\ ax7: p3 \sqsubseteq p6, ax8: p4 \sqsubseteq p7, ax9: p5 \sqsubseteq \forall s.p8, ax10: p6 \sqsubseteq \exists s.\neg p8\} \quad (2.1.1.1.1)$$

$$\dots \\ \text{Or}(X) = \text{true for } X = ax2, ax3, ax4, ax5, ax7, ax8, ax9, p7 \sqsubseteq p3; \quad (2.1.1.3.1)$$

$$\text{Or}(X) = \text{false for } X = ax1, ax6, ax10, p7 \sqsubseteq p5, p3 \sqsubseteq p8 \quad (2.1.1.4.1)$$

Remark 2.1.1.1.0: We test (T>((true for X)&(false for X))) as Eqs. 2.1.1.1.1 > ((2.1.1.3.1)&(2.1.1.4.1)). (2.1.1.5.1)

LET $p, q, r, s, t, u, v, w, x, y, z:$
 $p1, p2, p3, p4, p5, p6, p7, p8, s, y, z.$

$$\begin{aligned}
& (((\sim(q < p) \& \sim(r < p)) \& \sim(\sim(s < p) \& \sim(s < q))) \& ((\sim(t < q) \& \sim(t < r)) \& (\sim(u < r) \& \sim(v < s)))) \& \\
& (\sim((\#x \& w) < t) \& \sim((\%x \& \sim w) < u))) > \\
& (((((\sim(r < p) \& \sim(\sim(s < p)) \& \sim(s < q) \& \sim(t < q))) \& ((\sim(u < r) \& \sim(v < s)) \& \sim((\#x \& w) < t) \& \\
& \sim(r < v)))) = (z = z)) \& \\
& (((\sim(q < p) \& \sim(t < r)) \& \sim((\%x \& \sim w) < u)) \& (\sim(t < v) \& \sim(w < r))) = (z @ z)) ; \\
& \quad \text{TCTC TCTC TTTC TTTC} \\
& \quad \text{TTTT TTTC TTTT TTTC} \\
& \quad \text{TTTT **TF**TT TTTT TT**TF**} \\
& \quad \text{TTTT TT**TF** TTTT TT**TF**} \\
& \quad \text{TTTT TTTT TTTC TTTC} \\
& \quad \text{TTTT TTTT TTTT TTTC} \\
& \quad \text{TTTT TTTT TTTT TT**TF** } \times 2 \\
& \quad \text{TTTT **TF**TT TTTT TT**TF** } \times 2 \\
& \quad \text{TTTT TT**TF** TTTT TT**TF** } \\
& \quad \text{TTTT TTTT TTTT TT**TF** } \times 4 \\
& \\
& \quad \text{TTTT TTTT TTTT TTTT } \times 2 \\
& \quad \text{TTTT **TF**TT TTTT TT**TF**} \\
& \quad \text{TTTT TT**TF** TTTT TT**TF**} \\
& \quad \text{TTTT TTTT TTTT TTTT } \times 2 \\
& \quad \text{TTTT TTTT TTTT TT**TF** } \times 2 \\
& \quad \text{TTTT TN**TF** TTTT TTTN } \times 2 \\
& \quad \text{TTTT TT**TF** TTTT TT**TF** } \\
& \quad \text{TTTT TTTT TTTT TTTN } \times 2 \\
& \quad \text{TTTT TTTT TTTT TT**TF** } \quad 107 \text{ steps} \quad (2.1.1.5.2)
\end{aligned}$$

Remark 2.1.2.1.0: We map the unique relations from repairs of R1, R2, R3, R4, R5 of

R1={p4⊆p5, p7⊆p3}, R2={p4⊆p5, p7⊆p3}, R3={p7⊆p3}, R4={p4⊆p5}, R5={p4⊆p5, p7⊆p3},
as p4⊆p5, p7⊆p3 (2.1.2.1.1)

$\sim(t < s) \& \sim(r < v) ;$ (2.1.2.1.2)

Remark 2.1.3.1.0: We test the CDP to imply the repairs:

Eqs. 2.1.1.5.1 implies 2.1.2.1.1. (2.1.3.1.1)

$$\begin{aligned}
& (((((\sim(q < p) \& \sim(r < p)) \& \sim(\sim(s < p) \& \sim(s < q))) \& ((\sim(t < q) \& \sim(t < r)) \& (\sim(u < r) \& \sim(v < s)))) \& \\
& (\sim((\#x \& w) < t) \& \sim((\%x \& \sim w) < u))) > (((((\sim(r < p) \& \sim(\sim(s < p)) \& \sim(s < q) \& \sim(t < q))) \& \\
& ((\sim(u < r) \& \sim(v < s)) \& \sim((\#x \& w) < t) \& \sim(r < v)))) = (z = z)) \& (((\sim(q < p) \& \sim(t < r)) \& \\
& \sim((\%x \& \sim w) < u)) \& (\sim(t < v) \& \sim(w < r))) = (z @ z)) > (\sim(t < s) \& \sim(r < v)) ;
\end{aligned}$$

TTTT **F**N**F**N TTTT **F**F**F**N
FF**F**F **F**F**F**N TTTT **F**F**F**N
TTTT **F**T**F**T TTTT **F**F**F**T
FF**F**F **F**F**F**T TTTT **F**F**F**T
TTTT TTTT TTTT TTTT } × 2
FF**F**F **F**F**F**F TTTT TTTT }
TTTT **F**T**F**T TTTT **F**F**F**T } × 2
FF**F**F **F**F**F**T TTTT **F**F**F**T }
TTTT TTTT TTTT TTTT } × 2
FF**F**F **F**F**F**F TTTT TTTT }

TTTT **F**F**F**F TTTT **F**F**F**F
FF**F**F **F**F**F**F TTTT **F**F**F**F
TTTT **F**T**F**T TTTT **F**F**F**T
FF**F**F **F**F**F**T TTTT **F**F**F**T

$$\begin{array}{l}
TTTT \ TTTT \ TTTT \ TTTT \ } \times 2 \\
FFFF \ **FFFF** \ TTTT \ TTTT \ } \\
TTTT \ **FCFC** \ TTTT \ **FFFC** \ } \times 2 \\
FFFF \ **FFFT** \ TTTT \ **FFFT** \ } \\
TTTT \ TTTT \ TTTT \ TTTT \ } \times 2 \\
FFFF \ **FFFF** \ TTTT \ TTTT \ } \quad 115 \text{ steps} \quad (2.1.3.1.2)
\end{array}$$

Eqs. 2.1.1.5.2 and 2.1.3.1.2 as rendered are *not* tautologous. This means the example given as 2.1.1.1 is *not* tautologous, the oracles in 2.1.1.3.1 and 2.1.1.4.1 are *not* truthful, and the repairs in 2.1.2.1.1 are incorrect. This refutes the conjecture of ontology engineering.