

Shorter refutation of Gödel's completeness theorem

© Copyright 2019 by Colin James III All rights reserved.

Abstract: The completeness theorem rendered as $(\forall x.R(x,x)) \rightarrow (\forall x \exists y.R(x,y))$ is *not* tautologous. The application of Isabelle/HOL to prove the same also is *not* tautologous, to invalidate that tool. These demonstrations form a *non* tautologous fragment of the universal logic $\forall\exists\forall$.

We assume the method and apparatus of Meth8/ $\forall\exists\forall$ with Tautology as the designated proof value, **F** as contradiction, **N** as truthity (non-contingency), and **C** as falsity (contingency). The 16-valued truth table is row-major and horizontal, or repeating fragments of 128-tables, sometimes with table counts, for more variables. (See ersatz-systems.com.)

LET \sim Not, \neg ; + Or, \vee , \cup , \sqcup ; - Not Or; & And, \wedge , \cap , \square , $;$; \ Not And;
 $>$ Imply, greater than, \rightarrow , \Rightarrow , \mapsto , $>$, \supset , \Rightarrow ; $<$ Not Imply, less than, \in , $<$, **C**, \neq , \neq , \ll , \leq ;
 $=$ Equivalent, \equiv , $:=$, \Leftrightarrow , \leftrightarrow , \triangleq , \approx , \cong ; @ Not Equivalent, \neq ;
 $\%$ possibility, for one or some, \exists , \diamond , **M**; # necessity, for every or all, \forall , \square , **L**;
 $(z=z)$ **T** as tautology, \top , ordinal 3; $(z@z)$ **F** as contradiction, \emptyset , Null, \perp , zero;
 $(\%z>\#z)$ **N** as non-contingency, Δ , ordinal 1; $(\%z<\#z)$ **C** as contingency, ∇ , ordinal 2;
 $\sim(y < x)$ ($x \leq y$), ($x \subseteq y$), ($x \sqsubseteq y$); $(A=B)$ $(A \sim B)$.
 Note for clarity, we usually distribute quantifiers onto each designated variable.

From: en.wikipedia.org/wiki/Gödel%27s_completeness_theorem

Gödel's completeness theorem is a fundamental theorem in mathematical logic that establishes a correspondence between semantic truth and syntactic provability in first-order logic. ... The formula

$$(\forall x.R(x,x)) \rightarrow (\forall x \exists y.R(x,y)) \tag{1.1}$$

holds in all structures. By Gödel's completeness result, it must hence have a natural deduction proof.

$$\text{LET } q, r, s: x, R, y. \tag{1.2}$$

$$(r \& (\#q \& \#q)) > (r \& (\#q \& \%s)); \text{TTTT TTCC TTTT TTTT}$$

Remark 1.2: Not distributing the quantifiers as $(\#q \& (r \& (q \& q))) > ((\#q \& \%s) \& (r \& (q \& s)))$ produces the same truth value table.

Eq. 1.2 as rendered is *not* tautologous, refuting Gödel's completeness theorem.

Remark 1.3: The following paper in using an assistant to prove Gödel's completeness theorem further refutes the Isabelle/HOL tool itself: Margetson, J. (2014). Proving the completeness theorem within Isabelle/HOL. isa-afp.org/entries/Completeness-paper.pdf.