

## Refutation of predicative collapse and arithmetical comprehension

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**Abstract:** The four equations evaluated are *not* tautologous, hence refuting predicative collapsing principles such as arithmetical comprehension. Therefore these form a tautologous fragment of the universal logic  $V\mathbb{L}4$ .

We assume the method and apparatus of Meth8/ $V\mathbb{L}4$  with Tautology as the designated proof value, **F** as contradiction, **N** as truthity (non-contingency), and **C** as falsity (contingency). The 16-valued truth table is row-major and horizontal, or repeating fragments of 128-tables, sometimes with table counts, for more variables. (See ersatz-systems.com.)

LET  $\sim$  Not,  $\neg$ ; + Or,  $\vee, \cup, \sqcup$ ; - Not Or; & And,  $\wedge, \cap, \square, ;$ ; \ Not And;  
 $>$  Imply, greater than,  $\rightarrow, \Rightarrow, \mapsto, >, \supset, \rightsquigarrow$ ;  $<$  Not Imply, less than,  $\in, <, \subset, \neq, \neq, \neq, \leq$ ;  
 $=$  Equivalent,  $\equiv, :=, \Leftrightarrow, \leftrightarrow, \triangleq, \approx, \simeq$ ; @ Not Equivalent,  $\neq$ ;  
 $\%$  possibility, for one or some,  $\exists, \diamond, M$ ; # necessity, for every or all,  $\forall, \square, L$ ;  
 $(z=z)$  **T** as tautology, **T**, ordinal 3;  $(z@z)$  **F** as contradiction,  $\emptyset, \text{Null}, \perp$ , zero;  
 $(\%z\>\#z)$  **N** as non-contingency,  $\Delta$ , ordinal 1;  $(\%z\<\#z)$  **C** as contingency,  $\nabla$ , ordinal 2;  
 $\sim(y < x)$   $(x \leq y)$ ,  $(x \subseteq y)$ ,  $(x \sqsubseteq y)$ ;  $(A=B)$   $(A\sim B)$ .  
 Note for clarity, we usually distribute quantifiers onto each designated variable.

From: Freund, A. (2019). Predicative collapsing principles. [arxiv.org/pdf/1906.07448.pdf](https://arxiv.org/pdf/1906.07448.pdf)

**Abstract.** We show that arithmetical transfinite recursion is equivalent to a suitable formalization of the following:

For every ordinal  $\alpha$  there exists an ordinal  $\beta$  such that  $1+\beta\cdot(\beta+\alpha)$  (ordinal arithmetic) admits an almost order preserving collapse into  $\beta$ . (1.1.1)

LET  $p, q: \alpha, \beta$ .

$$(\#p\>\%q)\>(((\%p\>\#p)+q)\&(q+p)) ; \quad \mathbf{FN\!T\!T} \quad \mathbf{FN\!T\!T} \quad \mathbf{FN\!T\!T} \quad \mathbf{FN\!T\!T} \quad (1.1.2)$$

$$(\#p\>\%q)\>(((\%p\>\#p)+q)\&(q+p)))\>q ; \quad \mathbf{TCT\!T} \quad \mathbf{TCT\!T} \quad \mathbf{TCT\!T} \quad \mathbf{TCT\!T} \quad (1.2.2)$$

Arithmetical comprehension is equivalent to a statement of the same form, with  $\beta\cdot\alpha$  at the place of  $\beta\cdot(\beta+\alpha)$ . ... (2.1.1)

$$((\%p\>\#p)+q)\&p ; \quad \mathbf{FN\!E\!T} \quad \mathbf{FN\!E\!T} \quad \mathbf{FN\!E\!T} \quad \mathbf{FN\!E\!T} \quad (2.1.2)$$

$$(((\%p\>\#p)+q)\&p)\>q ; \quad \mathbf{TCT\!T} \quad \mathbf{TCT\!T} \quad \mathbf{TCT\!T} \quad \mathbf{TCT\!T} \quad (2.2.2)$$

**Remark 2.2.2:** We see the logical equivalence of Eqs. 1.2.2 and 2.2.2, meaning  $((1+\beta\cdot(\beta+\alpha))\rightarrow\beta)=((\beta\cdot\alpha)\rightarrow\beta)$ , which is probably not what the author intended.

The four equations evaluated are *not* tautologous, hence refuting predicative collapsing principles such as arithmetical comprehension.