A proof of Sophie Germain primes conjecture

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Abstract

In this paper a proof of the existence of an infinite number of Sophie Germain primes, is going to be presented. In order to do that, we analyse the basic formula for prime numbers and decide when this formula would produce a Sophie Germain prime, and when not. Originally very difficult problem (in observational space) has been transformed into a simpler one (in generative space) that can be solved by elementary math.

1 Introduction

A prime p is a Sophie Germain prime if 2p + 1 is prime, too [1]. In that case the prime number 2p + 1 is called safe prime. These special primes have applications in public key cryptography, pseudorandom number generation, and primality testing; see, for example, [2, 5, 7]. Originally, they have been used also in the investigation of cases of Fermat's last theorem [3]. It has been conjectured that there exist infinitely many Sophie Germain primes, but this was unproven (see for instance, [6]).

In this paper it is going to be proved that exists an infinite number of Sophie Germain primes. The problem is addressed in generative space, which means that prime numbers are not going to be analysed directly, but rather their representatives, in the other space, that can be used to produce them.

In a very similar manner, it is possible to prove Twin Prime Conjecture and Polignac's conjecture for cousin primes[4].

Remark: Prime numbers 2 and 3 are in a sense special primes, since they do not share some of the common features of all other prime numbers. For instance, every prime number, apart from 2 and 3, can be expressed in the form 6l + 1 or 6l - 1, where $l \in N$. So, in this paper most of the time we analyse prime numbers bigger than 3. It has to be said that both 2 and 3 are Sophie Germain primes, but that has no impact on the conclusion of this paper.

2 Proof

It is easy to check that any prime number (apart form 2 and 3) can be expressed in the form 6l + 1 or 6s - 1 ($l, s \in N$). Having that in mind it is easy to check that numbers in the form 6l + 1, could never be Sophie Germain primes since the safe prime is in the form

$$2(6l+1) + 1 = 12l + 3 = 3(4l+1)$$

and that is composite number divisible by 3. So, the prime number that can potentially be Sophie Germain prime must be in the form 6s - 1 and then the safe prime is going to be in the form 6(2s) - 1.

We denote any composite number (that is represented as a product of prime numbers bigger than 3) with CPN5. Also, we mark with mpl a number in the form 6l + 1, and with mps a number in the form 6s - 1 ($l, s \in N$). In that case, it is easy to check that any composite number CPN5 can be expressed in the form $mpl \times mpl$, $mps \times mps$ or $mpl \times mps$. So, if we have a number in the form 6k - 1 that is composite number it must hold

$$k = \frac{CPN5 + 1}{6}.$$
 (2.1)

Since CPN5 should be in the mps form, CPN5 can be generally expressed as a product $mpl \times mps$, or

$$mpl = 6x + 1$$
 and $mps = 6y - 1(x, y \in N)$,

which leads to

$$CPN5 = mpl \times mps = 6(6xy - x + y) - 1,$$
 (2.2)

or, due to symmetry

$$mpl = 6y + 1$$
 and $mps = 6x - 1$,

which leads to

$$CPN5 = mpl \times mps = 6(6xy + x - y) - 1.$$
 (2.3)

So, if we replace (2.2, 2.3) in (2.1) we obtain forms for all k that potentially cannot produce a prime number. Those forms are expressed by the following equation

$$k = \begin{cases} (6x - 1)y + x \\ (6x + 1)y - x \end{cases},$$
(2.4)

where $x, y \in N$.

Also, we know that safe prime which is generated by 2k cannot be composite. If the safe number is composite the following equation must hold

$$k = \frac{CPN5 + 1}{6 \times 2},$$
 (2.5)

where CPN5 is composite number in the *mps* form. Using the same analysis as in the previous case, and replacing (2.2) and (2.3) in (2.5), we obtain additional cases in which k cannot be used to produce Sophie Germain prime pairs, and they are defined by the following equation

$$k = \begin{cases} (6x-1)y + \frac{x}{2}, x \text{ is even} \\ (6x+1)y - \frac{x}{2}, x \text{ is even} \\ (6x-1)y + 3x + \frac{x-1}{2}, x \text{ is odd} \\ (6x+1)y - 3x - \frac{x+1}{2}, x \text{ is odd} \end{cases},$$
(2.6)

where $x, y \in N$. Equations (2.4) and (2.6) give a sufficient and necessary condition for k, so that it cannot be used for generation of the prime pairs in the form (p, 2p + 1). In order to prove that there are infinitely many prime pairs in the form (p, 2p + 1) we need to prove that exists infinitely many k that cannot be expressed in the form (2.4) or (2.6). In order to do it, the number of k that cannot be expressed in the forms (2.4, 2.6) is going to be calculated. First, we will check the form of (2.4, 2.6) for some values of x.

$\boxed{\text{Case } x = 1}$	$\underline{\text{Case } x = 2}$	$\underline{\text{Case } x = 3}$	$\underline{\text{Case } x = 4}$
k = 5y + 1	k = 11y + 2	k = 17y + 3	k = 23y + 4
k = 5y + 3	k = 11y + 1	k = 17y + 10	k = 23y + 2
k = 7y - 1	k = 13y - 2	k = 19y - 3	k = 5(5y - 1) + 1
k = 7y - 4	k = 13y - 1	k = 19y - 11	k = 5(5y - 1) + 3
$\underline{\text{Case } x = 5}$	$\underline{\text{Case } x = 6}$	$\underline{\text{Case } x = 7}$	$\underline{\text{Case } x = 8}$
k = 29y + 5	k = 7(5y+1) - 1	k = 41y + 7	k = 47y + 8
k = 29y + 17	k = 5(7y) + 3	k = 41y + 24	k = 47y + 4
k = 31y - 5	k = 37y - 6	k = 43y - 7	k = 7(7y - 1) - 1
k = 31y - 18	k = 37y - 3	k = 43y - 25	k = 7(7y - 1) + 3

From examples, we can see that if (6x - 1) or (6x + 1) represent composite number, k that is represented by that number has also representation by one of the prime factors of that composite number. This can be easily proved in the general case, by direct calculation using representations similar to (2.2, 2.3). Here only one case is going to be analysed. All other cases can be analysed analogously. In this case we assume

$$(6x - 1) = (6l + 1)(6s - 1),$$

where $(l, s \in N)$. From previous equation x can be expressed as

$$x = 6ls - l + s.$$

Having that in mind, and selecting one representation of k that includes form (6x - 1), we

have

$$k = (6x - 1)y - x = (6l + 1)(6s - 1)y - 6ls + l - s$$

or

$$k = (6l+1)(6s-1)y - s(6l+1) + l = (6l+1)((6s-1)y - s) + l,$$

which means

$$k = (6l+1)f + l$$

and that represents already existing form of the representation of k for factor (6l+1), where

$$f = (6s - 1)y - s.$$

So, we can see that all patterns for k that potentially result in composite number, include prime numbers and it is going to be checked is the number of k that cannot be represented by the models (2.4, 2.6) finite or infinite. In order to do it, a method similar to the sieve of Eratosthenes [8] is going to be used. When all numbers that can be represented in form

$$5y + 1$$
 and $5y + 3$.

are removed from natural numbers, it can be seen that $r_1 = 2/5$ of all natural numbers are removed. So, $c_1 = 1 - 2/5 = 3/5$ of all natural numbers cannot be represented by those two patterns and they still contain some k that could be used for representation of Sophie Germain primes. If, now, in addition, the natural numbers in the form

$$7y - 1$$
 and $7y - 4$,

are removed, then the ratio of removed numbers can be calculated by the following equation (together with previously removed numbers, and taking care that every removed number is calculated only once; basically, we apply the formula for calculation of the probability of occurring of two events that are not mutually exclusive $P(A \cup B) = P(A) + P(B) - P(A \cap B)$)

$$r_2 = r_1 + \frac{2}{7} - \frac{2}{7} \times r_1 = r_1 + \frac{2}{7}(1 - r_1) = \frac{2}{5} + \frac{2}{7}\left(1 - \frac{2}{5}\right) = \frac{20}{5 \times 7}$$

So, we have

$$c_2 = 1 - r_2 = 1 - r_1 - \frac{2}{7}(1 - r_1) = \left(1 - \frac{2}{7}\right) \times c_1 = \frac{3 \times 5}{5 \times 7},$$

of all natural numbers that potentially can be used for "generation" of Sophie Germain primes.

Now, we denote prime numbers bigger than 3 as p5, and define set S as

$$S = \{p5(1), p5(2), \dots, p5(i)\},\$$

where *i* is a natural number (in this moment we consider *i* bigger than 2), and p5(1) = 5, p5(2) = 7 and so on. So, after removal of *i*th pair of numbers related to the *i*th *p*5, we will have

$$r_i = 2^1 \sum_{p \in S} \frac{1}{p} - 2^2 \sum_{\substack{p,q \in S, \\ p \neq q}} \frac{1}{p \times q} + 2^3 \sum_{\substack{p,q,r \in S, \\ p \neq q \neq r}} \frac{1}{p \times q \times r} + \dots + (-1)^{i+1} 2^i \prod_{p \in S} \frac{1}{p},$$

and

$$c_i = 1 - r_i = \frac{\prod_{p \in S} (p-2)}{\prod_{p \in S} p}$$

This formula can be proved by using mathematical induction method (but also by some other methods). Here we briefly present the induction method. We have already seen that basis of the induction method for i = 2 is already satisfied. We suppose that formula holds for i = n. Now, we are going to prove that it holds for i = n + 1, too. So, since it holds

$$c_n = 1 - r_n = \frac{\prod_{j=1}^n (p5(j) - 2)}{\prod_{j=1}^n p5(j)}$$

also holds

$$r_n = 1 - c_n = \frac{\prod_{j=1}^n p5(j) - \prod_{j=1}^n (p5(j) - 2)}{\prod_{j=1}^n p5(j)}$$

When we remove the next pair of numbers that correspond to p5(n+1) we have

$$r_{n+1} = r_n + \frac{2}{p5(n+1)} - \frac{2}{p5(n+1)} \times r_n.$$

After a few elementary calculations, we obtain the following equation

$$r_{n+1} = \frac{\prod_{j=1}^{n+1} p5(j) - \prod_{j=1}^{n+1} (p5(j) - 2)}{\prod_{j=1}^{n+1} p5(j)} = 1 - \frac{\prod_{j=1}^{n+1} (p5(j) - 2)}{\prod_{j=1}^{n+1} p5(j)},$$

which means that c_{n+1} is defined by the following equation

$$c_{n+1} = \left(1 - \frac{2}{p5(n+1)}\right) \times c_n = \frac{\prod_{j=1}^{n+1} (p5(j) - 2)}{\prod_{j=1}^{n+1} p5(j)},$$
(2.7)

and that concludes the proof.

If we continue process until we remove all possible patterns (defined by (2.4, 2.6)) related to all prime numbers bigger than 3 (and that is an infinite number), the nominator c_n is given by

$$C = \lim_{n \to +\infty} \left(c_n \times \prod_{j=1}^n p5(j) \right) = \prod_{j=1}^{+\infty} (p5(j) - 2).$$

It can be easily concluded that C is an infinite number. Beside that, if we consider the case of irreducible c_n , the value of the nominator of irreducible c_n , at the end of the process, is given by the following equation

$$Cir = \lim_{n \to +\infty} (nominator(irreducible(c_n))) = \prod_{j=1}^{+\infty} (p5_n t(j) - 2),$$
(2.8)

where $p5_nt$ represents p5 numbers that do not have a twin prime bigger than them. Again, it can be easily concluded that *Cir* is an infinite number. That implies that the number of kthat can not be represented by (2.4, 2.6) is infinite. Why is it so, it is going to be explained in the next section, in an elementary manner. That concludes the proof that number of Sophie Germain primes is infinite. Here we are going to present an alternative way. From (2.7) we know that the following relation holds

$$c_{n+1} = \left(1 - \frac{2}{p5(n+1)}\right) \times c_n.$$
 (2.9)

That can be interpreted that in every step we remove 2/p5(n+1) numbers of what is left for potential representation of Sophie Germain primes. If we consider one additional process CA where we in every step remove 2/5 of the numbers that is left, we can represent it as

$$ca_{n+1} = \left(1 - \frac{2}{5}\right) \times ca_n. \tag{2.10}$$

From (2.9) and (2.10) we can easily conclude that the following equation holds

$$ca_{n+1} < c_{n+1}, n \in N.$$
 (2.11)

However, it is not difficult to be seen that removal of the 2/5 of the numbers from those that are left, will result in infinite numbers that cannot be removed, at the end of the process. In other words, infinite number of the natural numbers cannot be removed by the process CA (this will also be analysed in the next section, in an elementary way, by box and ball experiment). Having that in mind, and equation (2.11), we can also conclude that the process proposed by (2.4, 2.6), will leave an infinite number of the numbers that cannot be represented by (2.4, 2.6). That concludes an alternative the proof that number of Sophie Germain primes is infinite.

3 Explanation

In this section it is going to be explained in an elementary manner, why the fact that $Cir \rightarrow +\infty$ means that number of Sophie Germain primes is infinite.

In order to do that, it is going to be assumed that an infinite number of numbered balls (with all natural numbers written on them only once), as well as, infinite number of boxes of infinite size, are available. Sieve elimination process (SEP) is going to be explained by usage of the balls and boxes. In other words, for every SEP a corresponding BB (balls and boxes) experiment is going to be constructed. Also, we define sieve elimination thread (SET) as a infinite series of numbers (e.g. $5y - 1, y \in N$). Fact that balls are numbered is not in any sense critical for reasoning and making decisions, but, sometimes, is helpful in the sense that it is easier to understand underlying processes. For instance, if we remove a SET $(5y-1, y \in N)$ - that means that we remove one fifth of all balls - we can (but we do not have to) remove the balls with numbers defined by SET - important thing is to remove one fifth of all balls. As it is going to be shown in the following example, if you stick to the numbers on the balls that are moved (rather than number of the balls that are moved), numbers can, even, create a small problem.

Example 1

Now, two experiments are going to be analysed.

<u>First experiment</u>: Imagine that balls with all natural numbers on them (written just once) are in source box (SB), that has the size of the number of natural numbers, and that we have another experimental box (EB) of proper size.

In a moment 1 minute before midnight, we are going to move balls with numbers 1 to 10 from SB to EB, and remove the ball with number 10 on it. In the 1/2 minute before midnight, balls marked 11-20 are transferred to EB, and ball with number 20 is removed from EB. We continue the process at the moments $1/2^n$, $n \in N$ minute before midnight put balls with numbers n * 10 + 1 to (n + 1) * 10 and remove ball with number (n + 1) * 10. The question is: What is the number of the balls in the EB at midnight? And the answer is obvious and you can say it is infinite.

<u>Second experiment</u>: Again, imagine that balls with all natural numbers on them (written just once) are in one box (source box SB - that has the size of the number of natural numbers) and that we have another empty box (experimental box - EB) of proper size.

In a moment 1 minute before midnight, we are going to move balls with numbers 1 to 10 from SB to EB, and remove the ball with number 1 on it. In the $\frac{1}{2}$ minute before midnight,

balls marked 11-20 are transferred to EB, and ball with number 2 is removed from EB. We continue the process at the moments $\frac{1}{2^n}$, $n \in N$ minute before midnight - put balls with numbers n*10+1 to (n+1)*10 and remove ball with number n+1. The question is: What is the number of the balls in the EB at midnight? And the answer seems obvious and you can say it is infinite. However, if you are asked to specify a ball with any specific number that is still in the EB, you would not be able to do it. The reason is quite obvious - for any number you choose you can specify the moment in time in which the ball with that number has been removed from the EB.

So, numbers on the balls are not relevant for the reasoning - in both experiments, that were previously analysed, the process was following - put 10 balls in the EB and then remove one - or very simplified put nine balls in EB in every step. If there is no collapse of elementary reasoning (CER), we can safely conclude that number of the balls in the EB at midnight is infinite. For instance if we can conclude that

$$1 + 2 + 3 + 4 + \dots = -\frac{1}{12},$$

we obviously have a CER problem. Standard summation of infinite series of natural numbers give us as a solution plus infinity, and in that case we have no CER problem.

Now, the following problem could be analysed: If we create a SEP using infinitely many SETs and create corresponding BB experiment, could we know is the number of balls in the EB at the end of BB experiment, finite or infinite? It is going to be shown that it is possible to answer that question under assumption that we do not suffer from CER. In every experiment, from now on, we will move all the balls form SB to a EBs of size 1. Then, every $x, x \in N$ boxes would be fused and some fraction of the balls is going to be removed from each EB. The x is selected in such a way that it ensures that natural number of the balls is removed from the EBs. If the ratio of the number of the balls in EB and size of the EB, after removal, can be reduced, it has to be reduced before next step. Why is it so, it is going to be explained in the second experiment of Example 2.

Example 2

<u>First experiment</u>: The following SEP and corresponding BB experiment are going to be analysed (at the beginning, the balls whose number is equal to the number of natural numbers are in SB, and natural number of empty EBs of size 1 is available):

STEP 1 - (1 minute before midnight). Move all balls from SB to infinite number of EB of size 1. Then, fuse every 2 EBs (that, now, have size equivalent to the size of 2 balls), and then remove one ball from each EB and put it back in the SB. So, in the moment 1 minute before midnight a half of the balls (also we can say a half of the natural numbers - again it is not important what are the numbers on the balls) are removed from the EBs and put to SB, and each EB contains 1 ball. So, using notation from Proof Section we have that ratio of removed balls is $r_1 = 1/2$ and ratio of the balls still contained in EBs $c_1 = 1/2$. Now, we have interpretation of the c_1 that its nominator represents the number of the balls in the each EB, whose size is represented by its denominator.

STEP 2 - (1/2 minute before midnight). Create an infinite number of EBs of size 4 (by fusing every two EB's of size 2) and remove one ball from each EB and put it back in SB.

Now we have that ratio of removed balls is $r_2 = 1/2 + 1/4 = 3/4$ and ratio of the balls still contained in EBs $c_2 = 1 - r_2 = 1/4$, and now we have interpretation of the c_2 that its nominator represents the number of the balls in the each EB, whose size is represented by its denominator.

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STEP N - $(1/2^{N-1} \text{ minute before midnight})$. Fuse every 2 EBs of size $2^{n-1}, n \in N$ and obtain the EBs of size $2^n, n \in N$ and then remove from each EB one ball and put it back in SB. In that case ratio of removed balls (or the balls contained by SB) is

$$r_n = \frac{2^n - 1}{2^n},$$

while the ratio of the balls contained by the all EBs is defined by

$$c_n = \frac{1}{2^n},$$

with interpretation that each EB of size given by denominator of c_n contains the number of balls defined by the nominator of c_n , which is 1. So, without CER, we can conclude that at the midnight, we have in EB one ball (or a finite number of balls). And this is in complete agreement with our intuition that we can remove all balls (except one) from the basket by halving the number in every step. At the end the size of EB is equivalent to the number of natural numbers.

Second experiment: In the following experiment we are going to check what is going to happen if we omit to remove balls in one step after fusion. Lets say that it happens in step f. In the step f - 1 we had that

$$r_{f-1} = \frac{2^{f-1} - 1}{2^{f-1}},$$

and

$$c_{f-1} = \frac{1}{2^{f-1}}.$$

So, if we skip removal of the balls from the fused boxes in the step f and continue the normal process (with the ball removal) in the step f + 1 we'll have the following situation

$$r_{f+1} = \frac{2^{f+1} - 3}{2^{f+1}},$$

and

$$c_{f+1} = \frac{3}{2^{f+1}}.$$

After step f + k we will have the following situation

$$r_{f+k} = \frac{2^{f+k} - 2^k - 1}{2^{f+k}},$$

and

$$c_{f+k} = \frac{2^k + 1}{2^{f+k}}.$$

In this case we can see that at midnight nominator of c_n is going toward infinity, so the number of the balls that will be contained in the EB at the end of the process is going to be infinite. And this clearly is in accordance with our knowledge, since we have kept $1/2^{f}$ balls in the EB's at the step f (that are never going to be removed by the rest of the process) and that number is infinite.

Example 3

Now, we are going to present another example. Imagine that our SEP process has corresponding BB experiment that fuses 5 EB's in every step. After that, in every step, $b/5, b \in \{1, 2, 3\}$ of the balls is removed from the EBs in the first experiment, and 4/5 of the balls in the second experiment. We are going to check what are the outcomes of those experiments.

First experiment:

STEP 1 - (1 minute before midnight). Move all balls from SB to EBs of size 1, and fuse every 5 EBs to obtain the EBs of size 5 with 5 balls inside. Then, remove the *b* balls from every EB and put it back to SB. So, in this moment nominator of c_1 is 5 - b > 1.

STEP 2 - (1/2 minute before midnight). Again fuse every 5 EBs to obtain the EBs of size 25 with $5 \times (5 - b)$ balls inside. Then, remove the b/5 of the balls from every EB and put it back to SB. So, in this moment nominator of c_2 is $(1 - b/5) \times 5 \times (5 - b) = (5 - b)^2$.

STEP N- $(1/2^{N-1}$ minute before midnight). Fuse every 5 EBs Then, remove the b/5 of the balls from every EB and put it back to SB. So, in this moment nominator of $c_N =$ $(1-b/5) \times 5 \times c_{N-1} = (5-b)^N$.

So, we can conclude that c_N is increasing function of time and at midnight, without suffering from CER, we can conclude that the number of the balls in the EB is going to be infinite.

Second experiment:

STEP 1 - (1 minute before midnight). Move all balls from SB to EBs of size 1, and fuse every 5 EBs to obtain the EBs of size 5 with 5 balls inside. Then, remove the 4 balls from every EB and put it back to SB. So, in this moment nominator of c_1 is 1.

STEP 2 -(1/2 minute before midnight). Again fuse every 5 EBs to obtain the EBs of size 25 with 5 balls inside. Then, remove the 4/5 of the balls from every EB and put it back to SB. So, in this moment nominator of c_2 is $(1 - 4/5) \times 5 \times 1 = 1$.

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STEP N- $(1/2^{N-1})$ minute before midnight). Fuse every 5 EBs Then, remove the 4/5 of the balls from every EB and put it back to SB. So, in this moment nominator of $c_N = (1-4/5) \times 5 * c_{N-1} = 1$.

So, we can conclude that c_N is constant function of time and at midnight, without suffering from CER, the number of the balls in the EB is going to be finite.

From previous examples we can see that without CER, if the corresponding BB experiment of a given SEP results with nominator of the c_n that is nondecreasing number that tends toward infinity, then the number of balls in the EB at the end of the process is infinite, which means that SEP will leave infinite number of natural numbers that do not fulfil the rules specified by all SETs. If the nominator of c_n tends toward finite number when the times moves toward midnight, than a BB experiment results with some finite number of balls in the EB at the end of the process, which means that only a finite number of the natural numbers do not fulfil the rules specified by all SETs.

Now, we are going to construct the corresponding BB experiment for the SEP defined by (2.4, 2.6). Every SET will remove some balls from the EBs, after fusion of EBs. At the beginning, the balls whose number is equal to the number of natural numbers are in SB, and empty EBs of size 1 is available (the number of EBs is equal to the number of the balls). The goal of the process is to check if the number of the balls in the EB at the end of the process is finite or infinite, and nothing more than that. Notice, that only important fact is the **number of the balls in the EB** - the numbers on the balls are irrelevant, so in this case we will assume that balls are not marked with numbers.

STEP 1 - (1 minute before midnight). Move all balls from SB to EBs and fuse every p5(1) = 5 EBs to obtain the EBs of size 5 with 5 balls inside. Then, remove the 2 balls from every EB and put it back to SB. So, in this moment nominator of c_1 is 3.

STEP 2 - (1/2 minute before midnight). Fuse every p5(2) = 7 EBs and establish EB's of size $p5(1) \times p5(2) = 35$. Remove from all EB's $(2/p5(2)) \times p5(2) \times c_1 = 2 \times c_1$ balls and move it back to SB. In this moment nominator of $c_2 = (1-2/p5(2)) \times p5(2) \times c_1 = (1-2/7) \times 7 \times 3 = 15$. Now, we can see that ratio c_2 has to be reduced from 15/35, to 3/7.

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STEP N - $(1/2^{N-1} \text{ minute before midnight})$. Fuse every p5(N) EBs together and remove from each EB $(2/p5(N)) \times p5(N) \times c_{N-1} = 2 \times c_{N-1}$ balls and move it back to SB. So, in this moment nominator of c_N is $\prod_{j=1}^{N} (p5(j) - 2)$. If it is necessary, ratio c_N has to be reduced before the process is continued. If we continue process until we remove all SETs we can see that at the end of the process (at midnight), the number of balls in the EB will be infinite (equal to Cir defined by (2.8), if we do not suffer from CER. Once again, that means that number of k that cannot be defined by (2.4) and (2.6) is infinite, and that completes the explanation of the conclusion from previous section, that the number of Sophie Germain primes is infinite.

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