

Refutation of bisimulation

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Abstract: We evaluate two definitions of bisimilarity, both *not* tautologous. That refutes bisimulation, along with its proof tools of coinduction and Howe's congruence method.

We assume the method and apparatus of Meth8/VL4 with Tautology as the designated proof value, **F** as contradiction, **N** as truthity (non-contingency), and **C** as falsity (contingency). The 16-valued truth table is row-major and horizontal, or repeating fragments of 128-tables, sometimes with table counts, for more variables. (See ersatz-systems.com.)

LET \sim Not, \neg ; + Or, \vee, \cup ; - Not Or; & And, \wedge, \cap ; \ Not And;
 $>$ Imply, greater than, $\rightarrow, \Rightarrow, \mapsto, \succ, \supset, \vdash, \models, \rightarrow$; $<$ Not Imply, less than, $\in, \prec, \subset, \neq, \neq, \leftarrow$;
 $=$ Equivalent, $\equiv, :=, \iff, \leftrightarrow, \triangleq, \approx, \cong$ @ Not Equivalent, \neq ;
 $\%$ possibility, for one or some, \exists, \diamond, M ; # necessity, for every or all, \forall, \square, L ;
 $(z=z)$ **T** as tautology, \top , ordinal 3; $(z@z)$ **F** as contradiction, \emptyset , Null, \perp , zero;
 $(\%z<\#z)$ **C** as contingency, Δ , ordinal 1; $(\%z>\#z)$ **N** as non-contingency, ∇ , ordinal 2;
 $\sim(y < x)$ ($x \leq y$), ($x \subseteq y$); $(A=B)$ ($A \sim B$).
 Note: For clarity we usually distribute quantifiers on each variable as designated.

From: Pitts, A. (2011). Howe's method for higher-order languages.
 In D. Sangorgi, J. Rutten (eds.), Advanced topics in bisimulation and coinduction,
 Cambridge tracts in theoretical computer science. no. 52, chapter 5, pages 197-232.
cl.cam.ac.uk/~amp12/papers/howmho/howmho.pdf andrew.pitts@cl.cam.ac.uk

5. Howe's method for higher-order languages

5.1 Introduction

[A]lthough it is usually easy to see that a bisimilarity \simeq satisfies

$$\forall Q. P(x) \simeq P'(x) \Rightarrow P(Q) \simeq P'(Q) \quad (5.1.1.1)$$

LET $p, q, r, s, t: P, Q, P', x, Q'$

$$((p\&s)=(r\&s)) > ((p\#q)=(r\#q)); \quad \text{TTC TTCT TTT TTT} \quad (5.1.1.2)$$

for compatibility of \simeq we have to establish the stronger property

$$\forall Q, Q'. P(x) \simeq P'(x) \wedge Q \simeq Q' \Rightarrow P(Q) \simeq P'(Q') \quad (5.1.2.1)$$

$$(((p\&s)=(r\&s)) \& (\#q=\#t)) > ((p\#q)=(r\#t)); \quad \text{TTC TTCT TTT TTT, TTT TTT TTT TTT} \quad (5.1.2.1)$$

This is often hard to prove directly from the definition of \simeq

Eqs. 5.1.1.2 and 5.1.2.2 as rendered are *not* tautologous. This refutes bisimilarity, along with its proof tools of coinduction and Howe's congruence method.