

Refutation of correctness from Pratt-Floyd-Hoare logic

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Abstract: We evaluate the Pratt-Floyd-Hoare logic aimed at correctness of computer programs. For Pratt, four equations are tested, and for Hoare one is tested. None are tautologous. This refutes the Pratt-Floyd-Hoare logic for correctness.

We assume the method and apparatus of Meth8/VL4 with Tautology as the designated *proof* value, **F** as contradiction, **N** as truthity (non-contingency), and **C** as falsity (contingency). The 16-valued truth table is row-major and horizontal, or repeating fragments of 128-tables, sometimes with table counts, for more variables. (See ersatz-systems.com.)

LET $p, q, r, s, t, u, v: p, q, a, b, p', q', R$
 \sim Not, \neg ; $+$ Or, \vee, \cup ; $-$ Not Or; $\&$ And, \wedge, \cap ; \setminus Not And;
 $>$ Imply, greater than, $\rightarrow, \mapsto, \succ, \supset, \vdash, \models, \Rightarrow$; $<$ Not Imply, less than, $\in, \prec, \subset, \neq, \neq, \Leftarrow$;
 $=$ Equivalent, $\equiv, :=, \iff, \leftrightarrow, \triangleq$ @ Not Equivalent, \neq ;
 $\%$ possibility, for one or some, \exists, \diamond, M ; $\#$ necessity, for every or all, \forall, \square, L ;
 $(z=z)$ **T** as tautology, \top , ordinal 3; $(z@z)$ **F** as contradiction, $\emptyset, \text{Null}, \perp$, zero;
 $(\%z\<\#z)$ **C** as contingency, Δ , ordinal 1; $(\%z\>\#z)$ **N** as non-contingency, ∇ , ordinal 2;
 $\sim(y < x)$ ($x \leq y$), ($x \subseteq y$); $(A=B)$ ($A \sim B$).

From: Pratt, V.R. (1976). Semantical considerations of Floyd-Hoare logic.
 17th IEEE Foundations of Computer Science Conference.
boole.stanford.edu/pub/semcon.pdf pratt@cs.stanford.edu

Pratt weakest antecedent and strongest consequent:

$$P\{a\}Q \equiv \neq(P \supset \{a\}Q) \tag{W.1}$$

$$((p\&r)\&q) \sim \sim(p \succ (r\&q)); \quad \mathbf{TFTF \ TFTF \ TFTF \ TFTF} \tag{W.2}$$

Remark W: The proposition for in Eq. W.2 is *not* tautologous.

Pratt axiom 2, in handwritten note on margin:

$$P\{a \cup b\}Q \vdash P\{a\}Q \tag{A2n.1}$$

$$((p\&(r+s))\&q) \succ ((p\&r)\&q); \quad \mathbf{TTTT \ TTTT \ TTTF \ TTTT} \tag{A2n.2}$$

Remark A2n: The axiom as a handwritten note is *not* tautologous and not equivalent to the axiom.

Pratt axiom 4:

$$Q\{P\}P \wedge Q \tag{A4.1}$$

$$(q\&p)\&(p\&q); \quad \mathbf{FFFT \ FFFT \ FFFT \ FFFT} \tag{A4.2}$$

Pratt axiom 4, alternate:

$$P \supset Q \{P\} Q \quad (\text{A4.alt.1})$$

$$p > ((q \& p) \& q); \quad \text{TFTT TFTT TFTT TFTT} \quad (\text{A4.alt.2})$$

Remark A4: The original and alternate axioms are *not* tautologous and not equivalent.

Hoare "Rules of consequence":

$$P \supset Q, Q \{a\} R \vdash P \{a\} R \quad (\text{HR.1.1})$$

$$\begin{aligned} ((p > q) \& ((q \& r) \& v)) > ((p \& r) \& v); \quad & \text{TTTT TTTT TTTT TTTT (4),} \\ & \text{TTTT TTFT TTTT TTFT (4)} \end{aligned} \quad (\text{HR.1.2})$$

Remark HR1: The rule of consequence is *not* tautologous.

Eqs. W2, A2.n.2, A4.2, A4.alt.2, HR.1.2 are *not* tautologous and contradict the respective alternates. This refutes the Pratt-Floyd-Hoare logic for correctness.