

## Refutation of De Finettian logics of indicative conditionals

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**Abstract:** We evaluate De Finettian logics which specifies four conditionals, as attributed to Aristotle, Boethius, Cooper-Cantwell, Jeffrey, and six axioms. None is tautologous. This refutes their use to justify connexive logic.

We assume the method and apparatus of Meth8/VL4 with Tautology as the designated *proof* value, **F** as contradiction, **N** as truthity (non-contingency), and **C** as falsity (contingency). For results, the 16-valued truth table is row-major and horizontal, or repeating fragments of 128-tables, sometimes with table counts, for more variables. Reproducible transcripts for results are available. (See ersatz-systems.com.)

LET  $p, q, r, s: A, B, C, D;$   
 $\sim$  Not,  $\neg$ ;  $+$  Or,  $\vee, \cup$ ;  $-$  Not Or;  $\&$  And,  $\wedge, \cap$ ;  $\setminus$  Not And;  
 $>$  Imply, greater than,  $\rightarrow, \mapsto, \succ, \supset, \vdash, \models$ ;  $<$  Not Imply, less than,  $\in, \prec, \subset$ ;  
 $=$  Equivalent,  $\equiv, :=, \iff, \leftrightarrow, \triangleq$   $@$  Not Equivalent,  $\neq$ ;  
 $\%$  possibility, for one or some,  $\exists, \diamond, M$ ;  $\#$  necessity, for every or all,  $\forall, \square, L$ ;  
 $(z=z)$  **T** as tautology,  $\top$ , ordinal 3;  $(z@z)$  **F** as contradiction,  $\emptyset$ , Null,  $\perp$ , zero;  
 $(\%z\#z)$  **C** as contingency,  $\Delta$ , ordinal 1;  $(\%z>\#z)$  **N** as non-contingency,  $\nabla$ , ordinal 2;  
 $\sim(y < x)$  ( $x \leq y$ ), ( $x \subseteq y$ );  $(A=B)$  ( $A \sim B$ ).

From: Égré, P.; Rossi, L.; Sprenger, J. (2019).

De Finettian logics of indicative conditionals. [arxiv.org/pdf/1901.10266.pdf](https://arxiv.org/pdf/1901.10266.pdf)  
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**Remark 0:** Because formulas in the paper are not labelled, equations are keyed to the respective section and number of consecutive appearance.

$$A \rightarrow B |_{\top\top} \neg(A \rightarrow \neg B) \quad (3.1.5.1)$$

$$(p > q) > \sim(p > \sim q); \quad \mathbf{FTFT \ FTFT \ FTFT \ FTFT} \quad (3.1.5.2)$$

$$\neg A \vee B |_{\top\top} \neg(\neg A \vee \neg B) \quad (3.1.6.1)$$

$$\sim((\sim p + q) > \sim(\sim p + \sim q)) = (p = p); \mathbf{TFTF \ TFTF \ TFTF \ TFTF} \quad (3.1.6.2)$$

$$A \rightarrow B |_{\text{ss}} A \wedge B \quad (3.2.1.1)$$

$$(p > q) > (p \& q); \quad \mathbf{FTFT \ FTFT \ FTFT \ FTFT} \quad (3.2.1.2)$$

$$A \rightarrow B |_{\text{ss}} B \rightarrow A \quad (3.2.2.1)$$

$$(p > q) > (q > p); \quad \mathbf{TTFT \ TTFT \ TTFT \ TTFT} \quad (3.2.2.2)$$

$$A \rightarrow B |_{\top\top} A \wedge B \quad (3.2.3.1)$$

$$\sim((p>q)>(p\&q))=(p=p) ; \quad \mathbf{TFTF \ TFTF \ TFTF \ TFTF} \quad (3.2.3.2)$$

$$A \rightarrow B \#_{TT} B \rightarrow A \quad (3.2.4.1)$$

$$\sim((p>q)>(q>p))=(p=p) ; \quad \mathbf{FFTF \ FFTF \ FFTF \ FFTF} \quad (3.2.4.2)$$

$$(A \rightarrow B) \rightarrow A \quad (4.2.1.1)$$

$$(p>q)>p ; \quad \mathbf{FTFT \ FTFT \ FTFT \ FTFT} \quad (4.2.1.2)$$

$$\text{Aristotle's thesis: } \neg(\neg A \rightarrow A) \quad (5.3.1.1)$$

$$\sim(\sim p > p) = (p = p) ; \quad \mathbf{TFTF \ TFTF \ TFTF \ TFTF} \quad (5.3.1.2)$$

$$\text{Boethius' thesis: } (A \rightarrow C) \rightarrow \neg(A \rightarrow \neg C) \quad (5.3.2.1)$$

$$(p>r)>\sim(p>\sim r) ; \quad \mathbf{FTFT \ FTFT \ FTFT \ FTFT} \quad (5.3.2.2)$$

$$\text{Holds for Cooper-Cantwell conditional : } \neg(A \rightarrow B) \equiv_m (A \rightarrow \neg B) \quad (5.3.3.1)$$

$$\sim(p>q) = (p>\sim q) ; \quad \mathbf{FTFT \ FTFT \ FTFT \ FTFT} \quad (5.3.3.2)$$

$$\text{Holds for any Jeffrey conditional: } A \rightarrow B \#_{TT} \neg B \rightarrow \neg A \quad (5.6.1.1)$$

$$\sim((p>q)>(\sim q > \sim p)) = (p = p) ; \quad \mathbf{FFFF \ FFFF \ FFFF \ FFFF} \quad (5.6.1.2)$$

None of the ten equations above as rendered is tautologous. This denies that use to justify connexive logic with which we dispensed previously elsewhere.