

An even number greater than 2 equals the sum of two prime numbers.

Abstract , simulation basiclogicof synchronization algorithm, reasoning judgment and hypothesis contradiction

Goldbach conjecture

[integer theory] $2N=P_1+P_2$

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The number (prime number), also called prime number, has infinite number. A natural number greater than 1 cannot be divisible by any other natural number except 1 and itself, in other words, that the number does not have any other factor except 1 and itself; otherwise, it is called a composite.

According to the basic arithmetic theorem, every integer larger than 1 is either a prime number itself or can be written as the product of a series of prime numbers, and if the order of these prime numbers in the product is not considered, the form written out is unique. The minimum prime number is 2

set up : (natural number , N) $N > 1$

even number = $2N$

odd number = $2N - 1$

$\therefore 2N \div 2 = N$ (Satisfying integer solution)

\therefore natural number $N > 1$ (There's no prime in an even number.)

refer to: Reference: Euclid's theory of infinity of prime numbers

On the contrary, if the natural number $2N$ is not equal to the prime number,

On the contrary, prime numbers can only appear in odd numbers.

Or Goldbach's guess is right. $2N = P_1 + P_2$

Or Goldbach's guess doesn't hold. $2N \neq P_1 + P_2$

What are the conditions for establishing and what are the conditions for not holding

suppose: $2N \neq P_1 + P_2$

Equal: $N \neq P$ ($P =$ Prime number)

$N =$ even number

N=odd number

$$(\text{even number: } N) - 1 = S_1$$

$$(\text{odd number: } N) - 2 = S_2$$

$$2N - S_1 (\text{N=even number}) = L_1$$

$$2N - S_2 (\text{N=odd number}) = L_1$$

L_1 (Either prime or compound.)

Hypothesis: L_1 prime number

$$2N - L_1 = (\text{Either prime or compound.})$$

$$2N - L_1 = (\text{Assuming that it is equal to; } A_1 \times B_1 \times C_1)$$

$2N - S_1 = L_1$ (suppose: $L_1 = \text{composite number}$) $= A_1 \times B_1 \times C_1$ (Simulation. The basic logic of arithmetic.)

Hypothesis : $L_1 = \text{prime number}$

$$\text{like that: } 2N - L_1 = S_1 (\text{suppose: } S_1 = \text{composite number}) = A_1 \times B_1 \times C_1$$

Extraction prime number、 A_1, B_1, C_1

$$\left\{ \begin{array}{l} 2N - A_1 = L_2 (\text{Integer decomposition}) = A_2^n \\ 2N - B_1 = H_2 (\text{Integer decomposition}) = B_2^n \\ 2N - C_1 = M_2 (\text{Integer decomposition}) = C_2 \end{array} \right.$$

L_2, H_2, M_2 、 (Either prime or compound.)

Hypothesis; L_2, H_2, M_2 、 It's a compound number、

Extraction prime number: A_2, B_2, C_2 、

$$\left\{ \begin{array}{l} 2N - A_2 = L_3 (\text{Integer decomposition}) = A_3^n \\ 2N - B_2 = H_3 (\text{Integer decomposition}) = B_3^n \\ 2N - C_2 = M_3 (\text{Integer decomposition}) = C_3^n \end{array} \right.$$

L_3, H_3, M_3 , (Either prime or compound.)

Hypothesis: L_3, H_3, M_3 , It's a compound number.

Extraction prime number: A_3, B_3, C_3 ,

$$\left\{ \begin{array}{l} 2N - A_3 = L_4(\text{Integer decomposition}) = A_4^n \\ 2N - B_3 = H_4(\text{Integer decomposition}) = B_4^n \\ 2N - C_3 = M_4(\text{Integer decomposition}) = C_4^n \end{array} \right.$$

L_4, H_4, M_4 , (Either prime or compound.)

Hypothesis: L_4, H_4, M_4 , It's a compound number.

Extraction prime number: A_4, B_4, C_4 ,

$$\left\{ \begin{array}{l} 2N - A_4 = L_5(\text{Integer decomposition}) = A_5^n \\ 2N - B_4 = H_5(\text{Integer decomposition}) = B_5^n \\ 2N - C_4 = M_5(\text{Integer decomposition}) = C_5^n \end{array} \right.$$

L_5, H_5, M_5 (Either prime or compound.)

Hypothesis: L_5, H_5, M_5 , It's a compound number.

Extraction prime number: A_5, B_5, C_5 ,

.....Analog arithmetic logic (WY1 式)

(一)、 Either arithmetic logic, finite loop (prime cycle)

The prime cycle stands for $2N$. Using this arithmetic logic, we judge that we are all compound numbers. ($2N \neq P_1 + P_2$)

Simulation: arithmetic, finite cycle hypothesis for all complex numbers

$$2N - A = B^b \text{ (Extraction prime number: } B \text{)}$$

$$2N - B = C^c \text{ (Extraction prime number: } C \text{)}$$

$$2N - C = A^a \text{ (Extraction prime number: } A \text{)}$$

(二)、 Either the arithmetic logic is infinitely noncyclic (arithmetic prime is not cyclic) (infinite increment of different prime numbers)

The prime number is not cyclic, which means $2N$. Using this arithmetic logic, it's all a compound number. ($2N \neq P_1 + P_2$)

The representation of an infinitely noncyclic prime number (infinitely increasing)

Any $2N$ condition is finite

$2N - P$ (Infinite number of different primes: $2 < P < 2N$)

Finite and infinite contradictions、 Vice versa by the above logic ($2N = P_1 + P_2$)

Hypothesis: ($2N \neq P_1 + P_2$)

extract (一) Finite arithmetic cycle、 ($2N \neq P_1 + P_2$)

And then do the analog arithmetic judgment.

set up: ($2N \neq P_1 + P_2$)

Here, E_2 、 F_2 、 G_2 、 It can be prime or compound.

$$\textcircled{2} \begin{cases} 2N - 2A_2 = 2E_2 \div 2(E_2, \text{ Integer decomposition}) = S_2^n \\ 2N - 2B_2 = 2F_2 \div 2(F_2, \text{ Integer decomposition}) = W_2^n \\ 2N - 2C_2 = 2G_2 \div 2(G_2, \text{ Integer decomposition}) = R_2^n \end{cases}$$

Extraction prime number: S_2, W_2, R_2

$$\textcircled{1} \begin{cases} 2N - S_2 = E_3(E_2, \text{ Integer decomposition}) = S_3^n \\ 2N - W_2 = F_3(F_2, \text{ Integer decomposition}) = W_3^n \\ 2N - R_2 = G_3(G_2, \text{ Integer decomposition}) = R_3^n \end{cases}$$

E_3 、 F_3 、 G_3 、 (Either prime or compound.)

Hypothesis ; E_3 、 F_3 、 G_3 、 It's a compound number.

Extraction prime number: S_3, W_3, R_3

$$\textcircled{2} \begin{cases} 2N - 2S_3 = 2E_4 \div 2(E_4, \text{ Integer decomposition}) = S_4^n \\ 2N - 2W_3 = 2F_4 \div 2(F_4, \text{ Integer decomposition}) = W_4^n \\ 2N - 2R_3 = 2G_4 \div 2(G_4, \text{ Integer decomposition}) = R_4^n \end{cases}$$

Here, E_4 、 F_4 、 G_4 、 It can be prime or compound.

Extraction prime number: S_4, W_4, R_4 、

$$\textcircled{1} \begin{cases} 2N - S_4 = E_5(E_5, \text{ Integer decomposition }) = S_5^n \\ 2N - W_4 = F_5(F_5, \text{ Integer decomposition }) = W_5^n \\ 2N - R_4 = G_5(G_5, \text{ Integer decomposition }) = R_5^n \end{cases}$$

E_5, F_5, G_5 , (Either prime or compound.)

Hypothesis ; E_5, F_5, G_5 , It's a compound number.

Extraction prime number: S_5, W_5, R_5 ,

$$\textcircled{2} \begin{cases} 2N - 2S_5 = 2E_6 \div 2(E_6, \text{ Integer decomposition }) = S_6^n \\ 2N - 2W_5 = 2F_6 \div 2(F_6, \text{ Integer decomposition }) = W_6^n \\ 2N - 2R_5 = 2G_6 \div 2(G_6, \text{ Integer decomposition }) = R_6^n \end{cases}$$

Here, E_6, F_6, G_6 , It can be prime or compound. Extraction prime number, $S_6, W_6,$

$R_6, \dots\dots$

Analog arithmetic basic logic (WY2 式)

(三) 、 Or, a finite arithmetic logic loop. (Prime number, cyclic logic)

Finite arithmetical cyclic logic, representing $2N$ that uses this arithmetic logic to judge that all are complex numbers

($2N \neq P1 + P2$)

(四) 、 Or, infinite arithmetic, no loop. (Infinite increase in different prime numbers P)

$2N$ With this arithmetic logic, it's all complex numbers. ($2N \neq P1 + P2$)

$2N - P$ (Infinite, different prime number, $P < 2N$)

The contradiction between finite and infinite is based on the above logic。 ($2N \neq P1 + P2$) The assumption is not true

Suppose: ($2N \neq P1 + P2$) Only choice. (一) 、 (三) Finite cyclic arithmetic logic。

Then, the two arithmetical logic hypothesis problems are combined into one problem, reasoning and judgment.

Either, (WY2) all prime numbers are the same as (WY1).

Either, the prime number (WY2) is incrementally different from (WY1).

suppose:

Let: (WY2's) (prime cycle) all prime numbers be the same as (WY1's) prime numbers.

Let (WY1) compound number arithmetic cyclic logic S column

$A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow F \rightarrow G \rightarrow H \rightarrow A$

Abstract: simulation, S item arithmetic loop simulation:

$$(1) 2N - A = B^b$$

$$(2) 2N - B = C^c$$

$$(3) 2N - C = D^d$$

$$(4) 2N - D = E^e$$

$$(5) 2N - E = F^f$$

$$(6) 2N - F = G^g$$

$$(7) 2N - G = H^h$$

$$(8) 2N - H = A^a$$

$$(1) 2N - A = B^b \dots\dots\dots$$

$2N$, Not equal to the sum of two prime numbers ($2N \neq P_1 + P_2$)

(WY1) Compound arithmetic cyclic logic reference (WY2)

$$(2N - A = B^b) \quad (2N - 2A = 2B^d)$$

$$2N - A = B^b \text{ 和 } 2N - 2A = 2B^d$$

$$A = B^b - 2B^d$$

$$A = B^d(B^{b-d} - 2)$$

$$2N - B^d(B^{b-d} - 2) = B^b$$

$$2N - B^b + 2B^d = B^b$$

$$N = 2B^b - 2B^d$$

B is not a prime factor of $2N$

$2N \div B < \text{Unsatisfied integer solution} >$

Hypothesis: equality holds (with fractional solution)

The contradiction between integer theory and hypothesis

Conversely

(WY2) ② A(8) → Infer (1) the prime number of (1) can not be obtained from B

($2N - 2A = 2B^d$ The assumption is not true

Arbitrary form , $2N - 2E = 2B^m$

$$2N - E^e = D$$

$$2N - C = (2N - E^e)^d$$

$$C = 2N - (2N - E^e)^d$$

$$2N - B = [2N - (2N - E^e)^d]^c$$

$$B = 2N - [2N - (2N - E^e)^d]^c$$

substitution

$$2N - 2B = 2E^m$$

$$2N - 2[2N - (2N - E^e)^d]^c = 2E^m$$

So there's only the relationship between 2N and E here.

$$2N - 2[2N - (2N - E^e)^d]^c = 2E^m$$

$$N - [2N - (2N - E^e)^d]^c = E^m$$

$$c \geq 2$$

$$[2N - (2N - E^e)^d]^c \text{ 【Must appear. } (2N)^c \text{】}$$

like that , $N - (2N)^c$

$$N(1 - 2^c N^{c-1})$$

$$N \text{ 【K】} - E \text{ 【S】} = 2E^m$$

And here we have the relationship between the existence of prime factors of N and

E

The prime factor of N does not include E

If the equation holds, K or S has a fractional solution.

It is assumed that there is a contradiction between the theory of fraction and the theory of propositional integers.

Conversely, the assumption that $(2N \neq P_1+P_2)$ (一), (三) finite arithmetic cyclic logic (prime cycle) exists a finite cycle hypothesis does not hold.

Then (1), (3) there exists an infinite assumption of the existence of two forms (infinitely increasing different prime numbers)

$2N$, There are contradictions between finite values and infinitely different prime numbers

Conversely, the above arithmetic logic $(2N \neq P_1+P_2)$ does not hold.

Then by the above arithmetic logic $(2N=P_1+P_2)$ (computer program is WY2)

哥德巴赫猜想

命题大于 2 的偶数等于两个素数之和

抽象，模拟同步算术基本逻辑，推理判断和假设矛盾

【整数理论】 $2N=P_1+P_2$

哥德巴赫猜想，命题大于 2 的偶数等于两个素数之和

那么自然数，素数，

命题可以我们就可以判断属于【整数理论】

数(prime number)又称素数，有无限个。一个大于 1 的自然数，除了 1 和它本身外，不能被其他自然数整除，换句话说就是该数除了 1 和它本身以外不再其他的因数;否则称为合数。

根据算术基本定理，每一个比 1 大的整数，要么本身是一个质数，要么可以写成一系列质数的乘积;而且如果不考虑这些质数在乘积中的顺序，那么写出来的形式是唯一的。最小的

质数是 2

设：(自然数 N) $N > 1$

偶数 = $2N$

$$\text{奇数} = 2N - 1$$

$$\because 2N \div 2 = N (\text{满足整数解})$$

$$\therefore \text{自然数 } N > 1 (\text{偶数里面没有素数})$$

$$\text{要么哥德巴赫猜想正确、} 2N = P_1 + P_2$$

$$\text{要么哥德巴赫猜想不成立、} 2N \neq P_1 + P_2$$

成立的条件是什么、不成立的条件是什么

$$\text{假设: } 2N \neq P_1 + P_2$$

$$\text{当 } N \neq P (P \text{ 等于任意素数})$$

$$N = \text{偶数}$$

$$N = \text{奇数}$$

$$(\text{偶数: } N) - 1 = S_1$$

$$(\text{奇数: } N) - 2 = S_2$$

$$2N - S_1 (N = \text{偶数}) = L_1$$

$$2N - S_2 (N = \text{奇数}) = L_1$$

L_1 (要么是素数、要么是复合数)

假设: L_1 素数

$$2N - L_1 = (\text{要么是素数、要么是复合数})$$

$$2N - L_1 = (\text{假设等于 } A_1 \times B_1 \times C_1)$$

$$2N - S_1 = L_1 (\text{假设: } L_1 = \text{复合数}) = A_1 \times B_1 \times C_1 (\text{模拟.算术基本逻辑.})$$

如果 $L_1 = \text{素数}$ 。

$$\text{那么 } 2N - L_1 = S_1 (\text{假设: } S_1 = \text{复合数}) = A_1 \times B_1 \times C_1$$

抽取素数、 A_1, B_1, C_1

$$\left\{ \begin{array}{l} 2N - A_1 = L_2 (\text{整数分解}) = A_2^n \\ 2N - B_1 = H_2 (\text{整数分解}) = B_2^n \end{array} \right.$$

$$2N - C_1 = M_2 (\text{整数分解}) = C_2$$

L_2, H_2, M_2 、(要么是素数、要么是复合数)

假设: L_2, H_2, M_2 、是复合数、

抽取素数: A_2 、 B_2 、 C_2 、

$$\left\{ \begin{array}{l} 2N - A_2 = L_3(\text{整数分解}) = A_3^n \\ 2N - B_2 = H_3(\text{整数分解}) = B_3^n \\ 2N - C_2 = M_3(\text{整数分解}) = C_3^n \end{array} \right.$$

L_3 、 H_3 、 M_3 、（要么是素数、要么是复合数）

假设: L_3 、 H_3 、 M_3 、是复合数

抽取素数: A_3 、 B_3 、 C_3 、

$$\left\{ \begin{array}{l} 2N - A_3 = L_4(\text{整数分解}) = A_4^n \\ 2N - B_3 = H_4(\text{整数分解}) = B_4^n \\ 2N - C_3 = M_4(\text{整数分解}) = C_4^n \end{array} \right.$$

L_4 、 H_4 、 M_4 、（要么是素数、要么是复合数）

假设: L_4 、 H_4 、 M_4 、是复合数

抽取素数: A_4 、 B_4 、 C_4 、

$$\left\{ \begin{array}{l} 2N - A_4 = L_5(\text{整数分解}) = A_5^n \\ 2N - B_4 = H_5(\text{整数分解}) = B_5^n \\ 2N - C_4 = M_5(\text{整数分解}) = C_5^n \end{array} \right.$$

L_5 、 H_5 、 M_5 （要么是素数、要么是复合数）

假设: L_5 、 H_5 、 M_5 、是复合数

抽取素数: A_5 、 B_5 、 C_5 、

.....模拟算术逻辑 (WY1 式)

(一)、 要么、算术逻辑有限循环 (素数循环)
 素数循环代表 $2N$ 用这个算术逻辑判断都是复合数 ($2N \neq P_1 + P_2$)
 模拟: 都是复合数, 算术有限循环假设 (算术逻辑)

$$2N - A = B^b \text{ (抽取素数: } B)$$

$$2N - B = C^c \text{ (抽取素数: } C)$$

$$2N - C = A^a \text{ (抽取素数: } A)$$

(二)、 要么、算术逻辑无限不循环 (无限增加不相同素数)
 不循环代表 $2N$ 用这个算术逻辑判断都是复合数 ($2N \neq P_1 + P_2$)
 不循环无穷的素数代表 (无限增加不相同素数)

任意 $2N$ 有限值

$$2N - P \text{ (无限不相同素数 } 2 < P < 2N)$$

有限和无限矛盾、反之由上面逻辑 ($2N = P_1 + P_2$)

假设 ($2N \neq P_1 + P_2$)

抽取 (一) 有限算术逻辑循环、 ($2N \neq P_1 + P_2$)

再进行模拟算术判断

设 ($2N \neq P_1 + P_2$)

这里 E_2 、 F_2 、 G_2 、可以是素数或者是复合数

$$\begin{cases} 2N - 2A_2 = 2E_2 \div 2(E_2 \text{ 因式分解}) = S_2^n \\ \textcircled{2} \begin{cases} 2N - 2B_2 = 2F_2 \div 2(F_2 \text{ 因式分解}) = W_2^n \\ 2N - 2C_2 = 2G_2 \div 2(G_2 \text{ 因式分解}) = R_2^n \end{cases} \end{cases}$$

假设 ; E_2 、 F_2 、 G_2 、是复合数

抽取素数: S_2 、 W_2 、 R_2

$$\begin{cases} 2N - S_2 = E_3(E_2 \text{ 整数分解}) = S_3^n \\ \textcircled{1} \begin{cases} 2N - W_2 = F_3(F_2 \text{ 整数分解}) = W_3^n \\ 2N - R_2 = G_3(G_2 \text{ 整数分解}) = R_3^n \end{cases} \end{cases}$$

E_3 、 F_3 、 G_3 、 (要么是素数、要么是复合数)

假设 ; E_3 、 F_3 、 G_3 、是复合数

抽取素数: S_3 、 W_3 、 R_3

$$\begin{cases} 2N - 2S_3 = 2E_4 \div 2(E_4 \text{ 因式分解}) = S_4^n \end{cases}$$

$$\textcircled{2} \quad 2N - 2W_3 = 2F_4 \div 2(F_4 \text{ 因式分解}) = W_4^n$$

$$2N - 2R_3 = 2G_4 \div 2(G_4 \text{ 因式分解}) = R_4^n$$

这里 E_4 、 F_4 、 G_4 、可以是素数或者复合数

抽取素数： S_4 、 W_4 、 R_4 、

$$\textcircled{1} \quad \left\{ \begin{array}{l} 2N - S_4 = E_5(E_5 \text{ 整数分解}) = S_5^n \\ 2N - W_4 = F_5(F_5 \text{ 整数分解}) = W_5^n \end{array} \right.$$

$$2N - R_4 = G_5(G_5 \text{ 整数分解}) = R_5^n$$

E_5 、 F_5 、 G_5 、（要么是素数、要么是复合数）

假设； E_5 、 F_5 、 G_5 、是复合数

抽取素数： S_5 、 W_5 、 R_5 、

$$\textcircled{2} \quad \left\{ \begin{array}{l} 2N - 2S_5 = 2E_6 \div 2(E_6 \text{ 因式分解}) = S_6^n \\ 2N - 2W_5 = 2F_6 \div 2(F_6 \text{ 因式分解}) = W_6^n \end{array} \right.$$

$$2N - 2R_5 = 2G_6 \div 2(G_6 \text{ 因式分解}) = R_6^n$$

这里 E_6 、 F_6 、 G_6 、可以是素数或者是复合数

抽取素因数 S_6 、 W_6 、 R_6 、

.....模拟算术计算逻辑（WY2 式）

（三）、要么、有限算术逻辑循环（素数循环）

有限算术循环逻辑，代表 $2N$ 用这个算术逻辑判断都是复合数

$$(2N \neq P_1 + P_2)$$

（四）、要么、无限算术不循环(无限增加不相同素数) $2N$ 用这个算术逻辑判断都是复合数

$$(2N \neq P_1 + P_2)$$

不循环无穷的素数代表（无限递增）

$$2N - P \quad (\text{无限不相同素数、} P < 2N)$$

有限与无限矛盾由上面逻辑（ $2N \neq P_1 + P_2$ ）不成立

反之（ $2N = P_1 + P_2$ ）

假设 $(2N \neq P_1 + P_2)$ 只能选择 (一)、(三) 有限循环假设 (素数循环)
将两个算术逻辑假设问题合并成一个问题, 推理判断。

要么、(WY2) 全部素数与 (WY1) 全部相同

要么、(WY2) 所得到素数与 (WY1) 存在递增不相同素数

设: (WY2 式) (素数循环) 所有素数、与 (WY1 式) 素数全部相同

设: (WY1) 复合数算术循环逻辑 S 项列

$A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow F \rightarrow G \rightarrow H \rightarrow A$

抽象: 模拟, S 列项算术循环模拟:

$$(1) 2N - A = B^b$$

$$(2) 2N - B = C^c$$

$$(3) 2N - C = D^d$$

$$(4) 2N - D = E^e$$

$$(5) 2N - E = F^f$$

$$(6) 2N - F = G^g$$

$$(7) 2N - G = H^h$$

$$(8) 2N - H = A^a$$

$$(1) 2N - A = B^b \dots\dots\dots$$

$2N$ 不等于两个素数和 $(2N \neq P_1 + P_2)$

(WY2) 复合数算术循环逻辑参照 (WY1)

$$(2N - A = B^b) \quad (2N - 2A = 2B^d)$$

$2N - A = B^b$ 和 $2N - 2A = 2B^d$

$$A = B^b - 2B^d$$

$$A = B^d (B^{b-d} - 2)$$

$$2N - B^d(B^{b-d} - 2) = B^b$$

$$2N - B^b + 2B^d = B^b$$

$$N = 2B^b - 2B^d$$

B 不是 2N 的素因数

$2N \div B <$ 不满足整数解 $>$

假设：等式成立（有分数解）

整数理论和假设矛盾

反之

(WY2) ② A(8) \rightarrow 推导(1)得不到 B 的素数

($2N - 2A = 2B^d$) 不成立

任意一式, $2N - 2E = 2B^m$

$$2N - E^e = D$$

$$2N - C = (2N - E^e)^d$$

$$C = 2N - (2N - E^e)^d$$

$$2N - B = [2N - (2N - E^e)^d]^c$$

$$B = 2N - [2N - (2N - E^e)^d]^c$$

代入

$$2N - 2B = 2E^m$$

$$2N - 2[2N - (2N - E^e)^d]^c = 2E^m$$

那么这里只存在关系式 2N 与 E 的关系式

$$2N - 2[2N - (2N - E^e)^d]^c = 2E^m$$

$$N - [2N - (2N - E^e)^d]^c = E^m$$

$$c \geq 2$$

$$[2N - (2N - E^e)^d]^c \text{ 【必定出现 } (2N)^c \text{】}$$

那么, $N - (2N)^C$

$N(1 - 2^C N^{C-1})$

$N [K] - E [S] = 2E^m$

这里就产生 N 与 E 存在素因数关系

而 N 的素因数不包含 E

如果等式成立 K 或 S 存在分数解

假设存在分数和命题整数论之间存在矛盾。

反之假设 ($2N \neq P_1 + P_2$) 只能选择 (一)、(三) 有限循环假设 (素数循环) 存在一式有限循环假设不成立

那么 (一)、(三) 里面两式存在一式为无穷假设 (无限递增不相同素数)

$2N$ 是有限值与无限不相同素数存在矛盾

反之由上面算术逻辑 ($2N \neq P_1 + P_2$) 不成立

那么由上面算术逻辑 ($2N = P_1 + P_2$)

Reference literature :

Abstract, simulates basic arithmetic logic, reasoning judgment and hypothesis contradiction.

Integer theory

Abstract hypothesis: finite number of prime numbers

From small to large in order of order P_1 、 P_2 、 P_3 、 \dots 、 P_n

Simulating basic arithmetic Logic: multiply from small to large

$P_1 \times P_2 \times P_3 \times \dots \times P_n = N$

Then , $N+1$

Is a prime or not a prime

Reasoning judgment :

If $N + 1$ is a composite number ,

set up : $W = P_1, P_2, P_3, \dots, P_n$ (Arbitrary prime number)

$(N+1) \div W$

$N \div W$ (Satisfying integer solution)

$1 \div W$ (Unsatisfied integer solution)

Propositional condition is integer theory

but , $1 \div W$ (Unsatisfied integer solution) , Belong to fraction.

Does not belong to integer theory

So $N + 1$ is a composite or a prime.

$N+1$ The factorized prime factor is certainly not assumed. $P_1, P_2, P_3, \dots, P_n,$

Inside

There are also other prime numbers in addition to the assumed finite number of prime numbers. So the original assumption is not true. That is, there are infinitely many prime numbers.

Note: this article belongs to Euclidean academic theory ,

But I need to quote Euclidean theory to prove my theory in mathematics.

So the Euclidean theory is analyzed and rewritten.

Welcome to comment on my article