

## Refutation of the modal logic GL<sub>2</sub>

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**Abstract:** We evaluate the modal logic GL<sub>2</sub> in two axioms and for satisfiability. One of the two axioms is *not* tautologous, and the five formulas for PSpace complete satisfiability are *not* tautologous. Hence GL<sub>2</sub> is refuted.

We assume the method and apparatus of Meth8/VL4 with Tautology as the designated *proof* value, **F** as contradiction, N as truthity (non-contingency), and C as falsity (contingency). The 16-valued truth table is row-major and horizontal. The 16-valued truth table is row-major and horizontal, or repeating fragments of 128-tables, sometimes with table counts, for more variables. (See ersatz-systems.com.)

LET ~ Not, ¬; + Or, ∨; - Not Or; & And, ∧; \ Not And;  
 > Imply, greater than, →; < Not Imply, less than, ⇐  
 = Equivalent, ≡; @ Not Equivalent, ≠;  
 % possibility, for one or some, ∃, ∅; # necessity, for every or all, ∀, □;  
 ~( y < x ) ( x ≤ y ), ( x ⊆ y ).

From: Gabelaia, D.; Gogoladzeb, K.; Jibladze, M.; Kuznetsov, E.; Marxa, M. (2018).  
 Modal logic of planar polygons. arxiv.org/pdf/1807.02868.pdf e.kuznetsov@freeuni.edu.ge

LET p, q, r, s: p, q, r, γ.

We also present a slightly more intuitive and concise axiomatization of PL<sub>2</sub> by the following two formulas:

$$(I) p \rightarrow \square[\neg p \rightarrow \square(p \rightarrow \square p)] \quad (4.2.1)$$

$$p \# (\sim p \# (p \# p)) ; \quad \text{TNTN TNTN TNTN TNTN} \quad (4.2.2)$$

$$(II) \square[(r \wedge q) \rightarrow \gamma] \rightarrow [(r \wedge q) \rightarrow \diamond(\neg(r \wedge q) \wedge \diamond p \wedge \diamond \square \neg p)] \quad (4.3.1)$$

$$\#((r \& q) \# s) \# ((r \& q) \# (\sim(r \& q) \# (p \# \sim p))) ; \quad \text{TTTT TTTT TTTT TTCC} \quad (4.3.2)$$

$$\text{Where } \gamma \text{ is the formula } \diamond \square(p \wedge q) \wedge \diamond \square(\neg p \wedge q) \wedge \diamond \square(p \wedge \neg q). \quad (4.4.1)$$

$$s = (\#(p \& q) \# (\sim p \& q) \# (p \& \sim q)) ; \quad \text{TTTT TTTT FFFF FFFF} \quad (4.4.2)$$

$$\text{Substituting Eq. 4.4.1 into 4.3.1:} \quad (4.5.1)$$

$$\#((r \& q) \# (\#(p \& q) \# (\sim p \& q) \# (p \& \sim q))) \# ((r \& q) \# (\sim(r \& q) \# (\#p \# \sim p))) ; \quad \text{TTTT TTTT TTTT TTTT} \quad (4.5.2)$$

While Eq. 4.5.2 is tautologous as an axiom, Eq. 4.2.2 is *not* tautologous as an axiom. This means the

slightly more intuitive and concise axiomatization of  $PL_2$  is refuted.

LET  $p, q, r, s, t: \varphi, \psi, r, m, e$ .

Theorem 5.1. The satisfiability problem of our logic is PSpace complete.  
Let  $C$  be the conjunction of these formulas:

$$r, m, e \text{ are disjoint and one of them holds at each world.} \quad (5.0.1.1)$$

$$\begin{aligned} &(r+s)+t; \\ &\mathbf{FFFF} \text{ TTTT TTTT TTTT, TTTT TTTT TTTT TTTT} \end{aligned} \quad (5.0.1.2)$$

$$r \rightarrow \diamond m \quad (5.0.2.1)$$

$$\begin{aligned} &r > \% t; \\ &\text{TTTT CCCC TTTT CCCC, TTTT TTTT TTTT TTTT} \end{aligned} \quad (5.0.2.2)$$

$$m \rightarrow \diamond e. \quad (5.0.3.1)$$

$$\begin{aligned} &s > \% t; \\ &\text{TTTT TTTT CCCC CCCC, TTTT TTTT TTTT TTTT} \end{aligned} \quad (5.0.3.2)$$

$$C = (r \rightarrow \diamond m) \& (m \rightarrow \diamond e). \quad (5.0.4.1)$$

$$\begin{aligned} &((r+s)+t) \& ((r > \% s) \& (s > \% t)); \\ &\mathbf{FFFF} \text{ CCCC CCCC CCCC, TTTT CCCC TTTT TTTT} \end{aligned} \quad (5.0.4.2)$$

$$\begin{aligned} &r \wedge \square C \wedge \diamond(m \vee e) \wedge \varphi \wedge \square((m \vee e) \rightarrow \psi) \\ &\text{is satisfiable in our logic.} \end{aligned} \quad (5.1.1.1)$$

$$\begin{aligned} &((r \& \#(((r+s)+t) \& ((r > \% s) \& (s > \% t)))) \& (\% (s > t) \& p)) \& \#((s+t) > q); \\ &\mathbf{FFFF FFFF FFFF FFFF, FFFF FFFF FFFF FFFN} \end{aligned} \quad (5.1.1.2)$$

Eq. 5.1.1.2 as rendered is *not* tautologous and not contradictory, differing from the state of contradiction by one value N. This means  $PL_2$  is *not* satisfiable as PSpace complete.