

## Tree-3-Cover Ratio of Graphs: Asymptotes and Areas

Paul August Winter\*

[\\*winter@ukzn.ac.za](mailto:winter@ukzn.ac.za)

### Abstract

The graph theoretical ratio, the tree-cover ratio, involving spanning trees of a graph  $G$ , and a 2-vertex covering (a minimum set  $S$  of vertices such that every edge (or path on 2 vertices) of  $G$  has at least vertex end in  $S$ ) of  $G$  has been researched. In this paper we introduce a ratio, called the *tree-3-covering ratio with respect to  $S$* , involving spanning trees and a 3-vertex covering (a minimum set  $S$  of vertices of  $G$  such that every path on 3 vertices has at least one vertex in  $S$ ) of graphs. We discuss the *asymptotic convergence* of this tree-3-cover ratio for classes of graphs, which may have application in ideal communication situations involving spanning trees and 3-vertex coverings of extreme networks. We show that this asymptote lies on the interval  $[0, \infty)$  with the dumbbell graph (a complete graph on  $n-1$  vertices appended to an end vertex) has tree-3-cover asymptotic convergence of  $1/e$ , identical to the convergence in the secretary problem, and the tree-cover asymptotic convergence of complete graphs. We also introduce the idea of a *tree-3-cover area* by integrating this tree-3-cover ratio.

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Key words: spanning trees of graphs, vertex cover, 3-vertex cover, ratios, social interaction, network communication, convergence, asymptotes.

## 1. Introduction

We shall use the graph theoretical notation of [8] where our graphs are simple and connected. The order of graphs will be  $n$  and size  $m$ .

### Spanning trees

The graph-theoretical concept of *spanning trees* can be found in many real world applications, especially in social networking scenarios. For example, research in [2] involves work on sexual networks in an American high school which suggest that sexual networking involving individuals at the school are characterized by long chains or “spanning trees”, implying that a large part of the school had sexual contact with each another.

### Vertex cover

The importance of minimum 2-vertex coverings of a graph  $G$ , i.e. a minimum set  $S$  of vertices such that every path of  $G$  on 2 vertices has at least one vertex in  $S$ , occurs often in real life applications involving (extreme) networks with a large number of nodes (see the parameterized Vertex Cover problem in [5] and [9]). The idea of a 3-vertex covering of a graph  $G$  was introduced in [10]- this involved the smallest set  $S$  of vertices such that every path of  $G$  on 3 vertices has at least one vertex in  $S$ . This allowed for the investigation of the effect of the “activation” of  $S$  on all other vertices on paths of length at most 2 connected to  $S$ .

### Ratios

Ratios, such as expanders, Raleigh quotient (see [1]), the central ratio of a graph (see [4]) and eigen-pair ratio of classes of graphs (see [14]), Independence and Hall ratios (see [7]), tree-cover ratio (see [13]),  $h$ -eigen formation ratio (see [17]),  $t$ -complete sequence ratio (see [15]), chromatic-cover ratio (see [11]), chromatic-complete difference ratio (see [12]) and the eigen-complete difference ratio (see [16]), have been investigated.

## Spanning trees and 3-vertex cover

In this paper we combine the two ideas of spanning trees and (minimum) 3-vertex cover to introduce the idea of a *tree-3-cover ratio* of a graph. The importance of large numbers of vertices, which occurs in (extreme) networks, allowed for the investigation *asymptotic convergent* of this tree-3-cover ratio for different classes of graphs. We found that this asymptote lies on the interval  $[0, \infty)$  with the dumbbell graph (the graph consisting of a complete graph on  $n-1$  vertices appended to an end vertex) having tree-3-cover asymptotic convergence of  $1/e$  identical to the secretary problem and the tree-2-cover asymptotic convergence of the complete graph (see [6] and [13]). The idea of *area* is also introduced which involves the Riemann integral of this tree-3-cover ratio.

This ratio  $\frac{|S|t(H(S))}{t(K_n)}$  involving spanning trees and 3-vertex cover  $S$  with its asymptotic property and area of classes of graphs is presented below:

### 1.1.1 Definition

A (*minimum*) 3-vertex cover of  $G$  is a smallest set of vertices of  $G$  such that every path on 3 vertices has at least one vertex in  $G$ . If  $u$  is a vertex in  $S$  and  $v$  a vertex not in  $S$  connected to  $u$ , we say that  $v$  is connected to  $S$  by a path of length at most 2.

### 1.1 .2 Definition

Let  $t(G)$  be the number of spanning trees of a connected graph  $G$  of order  $n$ . Let  $S$  be a set of vertices of a *minimum 3-vertex cover* of  $G$ , and  $H(S)$  the subgraph of  $G$  induced by  $S$ . We consider only the 2 cases (i) *Either*  $H(S)$  is connected or (ii)  $H(S)$  is disconnected and consists of trees as components. In case (ii)  $t(H(S))$  is defined as  $t(H(S)) = 1$ .

Then the ratio:

$tc(G)_3^s = \frac{|S|t(H(S))}{t(G)}$  is the *tree-3-cover ratio* of  $G$  with respect to  $S$ .

Note: If  $H(S)$  is disconnected, and not trees as components, then one can consider spanning forests involving the components of  $H(S)$ , but such cases are not considered in this paper.

### 1.1.2 Definition

The importance of graphs with a large number of vertices is well known. If  $\langle$  is a class of graphs and  $tc(G)_3^s = \frac{|S|t(H(S))}{t(G)} = f(n)$  for each  $G \in \langle$ , where  $n$  is the order of  $G$ , then the horizontal asymptote of  $f(n)$  is denoted by:

$$tcasymp(\langle)_3^s = \lim_{n \rightarrow \infty} f(n)$$

This asymptote is called the *tree-3-cover asymptote* of  $\langle$  which is an indication of the behavior of the tree cover ratio when the graph has a large number of vertices, such as in extreme networks.

### **An ideal communication problem and tree-cover asymptote**

In [9] the communication problem is to select a minimal set  $S$  of placed sensor devices in a service area so that the all the nodes of service area is accessible by the minimal set of sensors. This can be adapted to a situation where there is a need for a minimal set  $S$  of placed sensor devices to communicate with all nodes that can be reached by paths of length at most 2 from  $S$ . Finding the minimal set of sensors can be modelled as a 3-vertex cover problem, where the 3-vertex cover set  $S$  facilitates the communications between the sensors and the nodes (on paths of length at most 2 from  $S$ ) of the service area in networks with a large number of nodes (vertices), i.e. in extreme networks. If  $H(S)$ , in the 3-tree cover definition, is connected, and  $M$  represents the vertices of  $G$  not in  $S$ , then each vertex of  $M$  is connected *by an path of length at most 2* (an *out-3-vertex path*) to vertex of  $H(S)$  which is part of a spanning tree. Thus the *ease* of communication between vertices of  $H(S)$  and  $M$  through the out-3-vertex paths, involving spanning trees, may be represented by this tree-3-cover ratio – the “ideal” case, involving large number of nodes, -which we believe is in the case of complete graphs. The more

difficult communication case may be in the situation involving paths, where this tree-cover asymptote is infinite.

## 2. EXAMPLES OF TREE-3-COVER RATIOS AND ASYMPTOTES

### 2.1 Complete graph

Let  $G$  be the complete graph  $K_n$  on  $n$  vertices.

Then a minimum 3-covering set of  $K_n$  is any subset of  $n-2$  vertices of  $K_n$ , and since  $t(K_n) = n^{n-2}$ ;  $t'(K_{n-2}) = (n-2)^{n-4}$  we have:

$$tc(K_n)_3^s = \frac{|S|t(K_{n-2})}{t(K_n)} = f(n) = \frac{(n-2)(n-2)^{n-4}}{n^{n-2}} = \frac{1}{(n-2)} \left( \frac{n-2}{n} \right)^{n-2} \text{ which behaves}$$

like  $\frac{1}{n}$  for  $n$  large, so that:

$$\Rightarrow tc_{asymp}(K_n)_3^s = \lim_{n \rightarrow \infty} f(n) = 0.$$

### 2.2 Cycles

The cycle  $C_n$  on  $n = 3k$  vertices has  $t(C_n) = n$ , and a minimum 3-vertex cover  $S$  will be the  $\frac{n}{3}$  vertices of the disconnected graph induced by every third vertex of the cycle, so that  $t(H(S)) = 1$  and  $|S| = \frac{n}{3}$ . Thus:

$$tc(C_n)_3^s = \frac{|S|t(H(S))}{t(C_n)} = f(n) = \frac{1}{3} \text{ so that}$$

$$tc_{asymp}C_n)_3^s = \frac{1}{3}.$$

### 2.3 Complete split-bipartite graph

Let  $K_{\frac{n}{2}, \frac{n}{2}}$  be the complete split-bipartite graph on  $n$  vertices.

Then  $t(K_{\frac{n}{2}, \frac{n}{2}}) = \binom{n}{2}^{n-2}$  and either partite set can be taken as a minimum 3-vertex cover  $S$  which yields  $t(H(S)) = 1$  so that

$$tc(K_{\frac{n}{2}, \frac{n}{2}})^s_3 = \frac{|S|t(H(S))}{t(K_{\frac{n}{2}, \frac{n}{2}})} = \frac{n}{2\binom{n}{2}^{n-2}} = \left(\frac{2}{n}\right)^{n-3} = f(n) \text{ so}$$

$$tc(K_{\frac{n}{2}, \frac{n}{2}})^s_3 = \left(\frac{2}{n}\right)^{n-3} \text{ and}$$

$$tcasymp(K_{\frac{n}{2}, \frac{n}{2}})^s_3 = 0.$$

## 2.4 Paths

Let  $P_n$  be a path on  $n = 3k$  of vertices. A minimum vertex cover  $S$  consists of every third vertex of  $P_n$ . Since  $|S| = \frac{n}{3}$ ,  $t(H(S)) = 1$  and  $t(P_n) = 1$  we have:

$$tc(P_n)^s_3 = \frac{|S|t(H(S))}{t(P_n)} = f(n) = \frac{n}{3} \text{ so that}$$

$$tcasymp(P_n)^s_3 = \infty$$

## 2.5 Wheel graph

The wheel graph  $W_n$  on  $n = 3k + 1$  vertices has a cycle of length  $3k$  with each vertex joined to a center. The number of spanning trees of this wheel is:

$$t(W_n) = \left(\frac{3+\sqrt{5}}{2}\right)^n + \left(\frac{3-\sqrt{5}}{2}\right)^n - 2 \text{ and the minimum vertex cover } S \text{ will involve every}$$

third vertex of the cycle and the center vertex. Thus:

$$t'(H(S)) = 1 \text{ and:}$$

$$tc(W_n)^s_3 = \frac{|S|t(H(S))}{t(W_n)} = f(n) = \frac{\frac{n-1}{3} + 1}{\left(\frac{3+\sqrt{5}}{2}\right)^n + \left(\frac{3-\sqrt{5}}{2}\right)^n - 2} \approx \frac{n}{6\left(\frac{3}{2}\right)^n}; \quad n \text{ large, so}$$

that:

$$tcasymp(W_n)^s_3 = 0$$

## 2.5 Ladder graph

The ladder graph  $L_{\frac{n}{2}, \frac{n}{2}}$  on an even number  $n$  of vertices has:

$$t\left(L_{\frac{n}{2}, \frac{n}{2}}\right) = \frac{(2+\sqrt{3})^{\frac{n}{2}} - (2-\sqrt{3})^{\frac{n}{2}}}{\sqrt{3}} \quad \text{and} \quad t(H(S)) = 1, \quad \text{where } S \text{ is taken as follows:}$$

Let  $P$  and  $P'$  be the two paths, each having  $\frac{n}{2}$  vertices, of the ladder, with edges between matched vertices of the two paths. Take  $S$  as the set of alternating vertices on  $P$  and  $P'$ , where the first vertex of  $P$  is selected and the second vertex of  $P'$  is selected, so that  $S$  will have  $\frac{n}{2}$  vertices. Then we have:

$$tc(L_{\frac{n}{2}, \frac{n}{2}})^s_3 = \frac{|S|t(H(S))}{t(L_{\frac{n}{2}, \frac{n}{2}})} = f(n) = \frac{n\sqrt{3}}{2(2+\sqrt{3})^{\frac{n}{2}} - 2(2-\sqrt{3})^{\frac{n}{2}}}.$$

Since  $(2+\sqrt{3})^{\frac{n}{2}}$  dominates  $(2-\sqrt{3})^{\frac{n}{2}}$  for large  $n$  we have:

$$f(n) = \frac{n\sqrt{3}}{2(2+\sqrt{3})^{\frac{n}{2}} - 2(2-\sqrt{3})^{\frac{n}{2}}} \approx \frac{n\sqrt{3}}{2(2+\sqrt{3})^{\frac{n}{2}}} \quad \text{for large } n \text{ so that:}$$

$$tcasympt(L_{\frac{n}{2}, \frac{n}{2}})^s_3 = 0$$

### 2.6 Star graph with rays of length 1

Let  $S_{n,1}$  be the star graph on  $n$  vertices with  $n-1$  rays of length 1. Then its centre is its minimum 3-covering set so that:

$$tc(S_{n,1})^s_3 = \frac{|S|t(H(S))}{t(S_{n,1})} = f(n) = 1. \text{ Hence:}$$

$$tcasympt(S_{n,1})^s_3 = 1$$

### 2.7 Star graph with $k$ rays of length 2.

Let  $S_{n,k(2)}$  be the star graph in  $n$  vertices with  $k$  rays of length 2 from its center so that  $n=2k+1$  (odd). The center is the minimum 3-vertex cover so that  $|S|=1$  and  $t(H(S))=1$  so that:

$$tc(S_{n,k(2)})^s_3 = \frac{|S|t(H(S))}{t(S_{n,k(2)})} = 1 \text{ and}$$

$$tcasympt(S_{n,k(2)})^s_3 = 1.$$

### 2.8 Sun graph

Take a cycle on  $\frac{n}{2}$  vertices,  $n = 4k$ , and attach an end vertex to each vertex of the cycle to form the sun graph  $SN_n$  on  $n$  vertices. Since  $t(SN_n) = n$  and  $S$  consists of every alternate vertex of the cycle so that  $t(H(S)) = 1; |S| = \frac{n}{4}$ . Hence:



$$tc(SN_n)^{s_3} = \frac{|S|t(H(S))}{t(SN_n)} = \frac{n1}{4n} = \frac{1}{4} \text{ so that:}$$

$$tc(SN_n)^{s_3} = \frac{1}{4} \text{ and } tcasympt(SN_n)^{s_3} = \frac{1}{4}.$$

## 2.8 Dumbbell graph

Let  $D_n^2$  be the dumbbell graph consisting of two disjoint copies, A and B, of  $K_{\frac{n}{2}}$  joined by an edge  $uv$ .

For each spanning tree of A we get  $\left(\frac{n}{2}\right)^{\frac{n}{2}-2}$  spanning trees of  $D_n^2$  through the edge  $uv$ . Thus:

$$t(D_n^2) = \left(\frac{n}{2}\right)^{\frac{n}{2}-2} \left(\frac{n}{2}\right)^{\frac{n}{2}-2} = \left(\frac{n}{2}\right)^{n-4}$$

A 3-vertex cover of A will consist of any set P of  $\frac{n}{2} - 2$  vertices of A containing  $u$ .

A 3-vertex cover of B will consist of any set Q of  $\frac{n}{2} - 2$  vertices of B containing  $v$ .

Since each spanning tree of a 3-covering of  $D_n^2$  must contain  $uv$ , the subgraph  $H(P \cup Q)$  induced by  $S = P \cup Q$  will contain the following number of spanning trees:

$$t(H(S)) = \left(\frac{n}{2} - 2\right)^{\frac{n}{2}-4} \left(\frac{n}{2} - 2\right)^{\frac{n}{2}-4} = \left(\frac{n}{2} - 2\right)^{n-8}$$

.

Thus:

$$tc(D_n^2)^{s_3} = \frac{|S|t(H(S))}{t(SN_n)} = \frac{(n-4)\left(\frac{n}{2}-2\right)^{n-8}}{\left(\frac{n}{2}\right)^{n-4}} = \frac{(n-4)^{n-7}}{2^{n-8} 2^{-n+4} n^{n-4}} = \frac{2^4}{(n-4)^3} \left(\frac{n-4}{n}\right)^{n-4}.$$

Thus:  $tcasympt(SN_n)^s_3 = 0$ .

## 2.9 Lollipop graph

Let  $LP_{n-1,1}$  be the lollipop graph consisting of a complete graph  $F$  on  $n-1$  vertices with vertex  $u$  joined to a single end vertex.

The number of spanning of  $LP_{n-1,1}$  will be  $(n-1)^{n-3}$ .

A 3-vertex cover of  $LP_{n-1,1}$  will consist of a set  $S$  of  $n-2$  vertices of  $F$  *not* including  $u$ .

Thus:

$t(H(S)) = (n-2)^{n-4}$  so that:

$$tc(LP_{n-1,1})^s_3 = \frac{|S|t(H(S))}{t(SN_n)} = \frac{(n-2)(n-2)^{n-4}}{(n-1)^{n-3}} = \left(\frac{n-2}{n-1}\right)^{n-3} = \left(1 - \frac{1}{n-1}\right)^{n-3}.$$

$$\text{let } y = \left(1 - \frac{1}{n-1}\right)^{n-3} \Rightarrow \ln y = (n-3) \ln\left(1 - \frac{1}{n-1}\right) = \frac{\ln\left(1 - \frac{1}{n-1}\right)}{\frac{1}{n-3}}.$$

Letting  $n$  go to infinity we get:

$$\lim_{n \rightarrow \infty} \ln y = \lim_{n \rightarrow \infty} \frac{\ln\left(1 - \frac{1}{n-1}\right)}{\frac{1}{n-3}} = \lim_{n \rightarrow \infty} \frac{\left(\frac{1}{1 - \frac{1}{n-1}}\right)(n-1)^{-2}}{-(n-3)^{-2}} = \lim_{n \rightarrow \infty} - \left(\frac{1}{1 - \frac{1}{n-1}}\right) \left(\frac{n-3}{n-1}\right)^2 = -1$$

Thus:  $\lim_{n \rightarrow \infty} y = e^{-1} = tcasympt(LP_{n-1,1})^s_3$ , identical to the secretary problem.

Theorem

The tree-cover ratios and tree-cover asymptotes of the following graphs are:

$$tc(K_n)^s_3 = \frac{1}{n-2} \left( \frac{n-2}{n} \right)^{n-2} \text{ and } tcasymp(K_n)^s_3 = 0.$$

$$tc(C_n)^s_3 = \frac{1}{3} \text{ and } tcasymp(C_n)^s_3 = \frac{1}{3}; n = 3k.$$

$$tc(K_{\frac{n}{2}, \frac{n}{2}})^s_3 = \left( \frac{2}{n} \right)^{n-3} \text{ and } tcasymp(K_{\frac{n}{2}, \frac{n}{2}})^s_3 = 0$$

$$tc(P_n)^s_3 = \frac{n}{3} \text{ and } tcasymp(P_n)^s_3 = \infty; n = 3k$$

$$tc(W_n)^s_3 = \frac{\frac{n-1}{3} + 1}{\left( \frac{3+\sqrt{5}}{2} \right)^n + \left( \frac{3-\sqrt{5}}{2} \right)^n - 2} \text{ and } tcasymp(W_n)^s_3 = 0; n = 3k + 1$$

$$tc(L_{\frac{n}{2}, \frac{n}{2}})^s_3 = \frac{n\sqrt{3}}{(2+\sqrt{3})^n - (2-\sqrt{3})^n} \text{ and } tcasymp(L_{\frac{n}{2}, \frac{n}{2}})^s_3 = 1$$

$$tc(S_{n,1})^s_3 = 1 \text{ and } tcasymp(S_{n,1})^s_3 = 1$$

$$tc(S_{n,k(2)})^s_3 = 1 \text{ and } tcasymp(S_{n,k(2)})^s_3 = 1$$

$$tc(SN_n)^s_4 = \frac{1}{4} \text{ and } tcasymp(SN_n)^s_3 = \frac{1}{4}; n = 4k.$$

$$tc(D_n^2)^s_3 = \frac{2^4}{(n-4)^3} \left( \frac{n-4}{n} \right)^{n-4} \text{ and } tcasymp(D_n^2)^s_3 = 0.$$

$$tc(LP_{n-1,1})^s_3 = \left( 1 - \frac{1}{n-1} \right)^{n-3} \text{ and } tcasymp(LP_{n-1,1})^s_3 = \frac{1}{e}.$$

Corollary

The tree-3-cover asymptote for all classes of graphs lies on the interval  $[0, \infty]$ .

### 3. TREE-COVER AREA OF CLASSES OF GRAPHS

We introduce another dimension by integrating this tree-cover ratio.

#### 3.1 Definition

If  $\kappa$  is a class of graphs and  $tc(G)^s_3 = \frac{|S|t(H(S))}{t(G)} = f(n)$  for each  $G \in \kappa$ , where  $n$  is the size of  $G$  and  $G$  has  $m$  edges, then the *tree cover area* of  $\kappa$  is defined as:

$$tcA^3_{\kappa(n)} = \frac{2m}{n} \int f(n)dn; \quad tcA^3_{\kappa(p)} = 0 \text{ for } \min p \text{ defined}$$

#### Average degree

The value  $\frac{2m}{n}$  represents the *average degree* of a graph  $G$ .

#### Tree-cover height

For complete graphs, the length of the longest path is  $(n-1)$  so that we refer to the integral part of the definition as the *tree-3-cover height* of the graph.

#### 3.1 Example- cycle

If  $C_n$  is a cycle on  $n = 3k$  vertices, then:

$$tc(C_n)^s_3 = \frac{|S|t(H(S))}{t(C_n)} = f(n) = \frac{1}{3} \text{ so that the tree-cover height of cycles is:}$$

$\int \frac{1}{3}dn$  which gives the tree-cover area of cycles as:

$$tcA_{C_n}^3 = \frac{2n}{n} \int \frac{1}{3} dn = 2\left(\frac{n}{3} + c\right); tcA_{C_3}^3 = 0 \Rightarrow c = -1$$

Theorem

$$tcA_{C_n}^3 = \frac{2}{3}n - 2; n = 3k.$$

### 3.2 Example- the path

If  $P_n$  is a path on  $n = 3k$  number of vertices then:

$$tc(P_n)^s = \frac{|S|t(H(S))}{t(P_n)} = f(n) = \frac{n}{3} \text{ so that:}$$

$$tcA_{P_n}^3 = \frac{(2n-2)}{n} \int \frac{n}{3} dn = \frac{2(n-1)}{n} \left(\frac{n^2}{3} + c\right); tcA_{P_3}^3 = 0 \Rightarrow c = -3$$

Theorem

$$tcA_{P_n}^3 = \frac{2(n-1)}{n} \left(\frac{n^2}{3} - 3\right); n = 3k.$$

### 3.3 Example- star graph with rays of length 1

$$tcA_{S_{n,1}}^3 = \frac{(2n-2)}{n} \int dn = \frac{(2n-2)}{n} (n+c); tcA_{S_{1,1}}^3 = 0 \Rightarrow c = -2$$

Theorem

$$tcA_{S_{n,1}}^3 = \frac{(2n-2)}{n} (n-2)$$

### 3.4 Example= star graph with rays of length 2

$$tcA_{Sk(2)}^3 = \frac{2n-2}{n} \int dn = \frac{(2n-2)}{n} (n+c); tcA_{S1(2)}^3 = 0 \Rightarrow c = -3$$

Theorem

$$tcA_{Sk(2)}^3 = \frac{(2n-2)}{n} (n-3)$$

### 3.5 Sun graph

$$tcA_{SNn}^3 = 2 \int \frac{1}{4} dn = 2 \left( \frac{n}{4} + c \right); n = 4k; tcA_{SN8}^3 = 0 \Rightarrow c = -2$$

Theorem

$$tcA_{SNn}^3 = \frac{n}{2} - 4.$$

## 4. CONCLUSION: KNOWN AND NEW RESULTS

### 4.1 Combining spanning trees and 3-vertex coverings

In this paper we combined the concepts of spanning trees  $t(G)$  and a minimum 3-vertex cover,  $S$ , of a graph  $G$ , to introduce a *new* concept of a tree-3-cover ratio of  $G$  (where  $H(S)$  is the induced subgraph of  $G$  induced by a minimum 3-vertex covering  $S$  of  $G$ ):

$$\frac{|S|t(H(S))}{t(G)}$$

This ratio was motivated by the possible importance of 3-vertex coverings in sensor activation, the tree-cover ratio of [13], and that the general tree-3-cover ratio for lollipop graphs, as a function of the order  $n$  of such graphs, is

$$\left( 1 - \frac{1}{n-1} \right)^{n-3}.$$

This ratio has the asymptotic convergence of  $1/e$ , which is identical to the probability of the best applicant being selected in the secretary problem. These considerations resulted in the investigation of the asymptotic convergence of the tree-3-cover ratio of classes of graphs. We introduced integration of the tree-3-cover ratio which allowed for the idea of tree-cover area of classes of graphs.

We propose that the tree-cover asymptote of the sun graph on  $n=4k$  vertices is the smallest amongst all such possible positive tree-3-cover asymptotes of classes of graphs. Future research may involve considering the tree-3-cover ratio of the complement of classes of graphs discussed here. We could have considered the reciprocal of the tree-cover ratio, i.e. the ratio:

$$(tc(G)_3)^{-1} = \frac{t(G)}{|S|t(H(S))}.$$

For example, the reciprocal of the tree-3-cover ratio of lollipop graphs would have the asymptotic convergence of  $e$ , while paths on  $3k$  number of vertices would have a reciprocal tree-cover asymptote of  $0$  (which is the same as the tree-cover asymptote of complete-split bipartite graphs) and (reciprocal) tree-cover area of

$$\frac{(2n-2)}{n} \int \frac{3}{n} dn = \frac{2(n-1)}{n} (3 \ln n + c).$$

#### 4.2 known and new results: ratios, asymptotes and areas

For the complete graph on  $n$  vertices the following are **known results**:

The *vertex expansion ratio*:  $\min_{|S| \leq \frac{n}{2}} \frac{|\partial(S)|}{|S|} = \frac{n/2}{n/2} = 1$  which has *asymptote* 1 (see [1])

The *Hall ratio*:  $\dots(G) = \max \left( \frac{|V(H)|}{r(H)} \right) = \frac{n}{1}$  which *converges* to infinity (see [7]).

The *integral eigen-ratio*, i.e the ratio of  $a+b$  to  $ab$ , where  $a$  and  $b$  and two, distinct non-zero eigenvalues whose sum and product is integral, is:

$\frac{n-2}{1-n}$  which converges to -1 and:

The *eigen-area*:  $(n-1)(n - \ln(n-1))$  (see [14]).

The *central radius ratio* is  $\frac{rad(G)}{n} = \frac{n}{n} = 1$  which has *asymptote* 1 (see [4]).

The *tree-cover ratio* (or tree-2-cover ratio) is  $tc(G)^s = \frac{|S|t(H(S))}{t(G)} = \left(1 - \frac{1}{n}\right)^{n-2}$  with *asymptote*  $1/e$  (see [13]).

The *H-eigen formation ratio* of the graph G, on m edges, with H-decomposition. Is:

$ratio_H E(G) = [E(G) - E^H(G)]/m$  so that for the complete graph we get:

$ratio_{K_2}(K_n) = \frac{2(-n^2 + 3n - 2)}{n(n-1)}$  with *asymptote* -2 (see[17]).

The *chromatic-cover ratio* is  $cov\{t^S(K_n)\} = \frac{|S|t(H(S))}{n^t(K_n)} = \frac{(n-1)^2}{n^2}$  with *asymptote* 1 (see [11]).

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