Can the external directed edges of a complete graph form a radially symmetric field at long distance?

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Abstract

Using a numerical method, the external directed edges of a complete graph are tested for their level of fitness in terms of how well they form a radially symmetric field at long distance (e.g., a test for the inverse square law in 3D space). It is found that the external directed edges of a complete graph can very nearly form a radially symmetric field at long distance if the number of graph vertices is great enough.

1 Introduction

Complete graphs have been used to construct a model of quantum gravity [1]. It is considered here that a complete graph G_1 consists of:

- 1. $n(G_1)$ vertices $V(G_1)$ that are uniformly distributed along a shell $S(G_1)$ of radius $r(G_1)$.
- 2. $(n(G_1)^2 n(G_1))/2$ internal non-directed edges $I(G_1)$ (e.g., line segments) that join the vertices together.
- 3. $n(G_1)^2 n(G_1)$ external directed edges $E(G_1)$ (e.g., rays) that are extensions of $I(G_1)$.

See Figure 1 for a diagram of a complete graph where $n(G_1) = 5$.

It seems fundamentally important to question whether or not the external directed edges $E(G_1)$ can form a radially symmetric field at long distance.

2 Method

If the field is to be considered radially symmetric, then the following two fitness criteria must be met:

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- 1. With regard to a second shell $S(G_2)$ of larger radius $r(G_2) > r(G_1)$, the $n(G_2) = n(G_1)^2 - n(G_1)$ vertices $V(G_2)$ corresponding to where the external directed edges $E(G_1)$ intersect with $S(G_2)$ should be uniformly distributed along $S(G_2)$.
- 2. The external directed edges $E(G_1)$ should be normal to $S(G_2)$ at their respective intersection vertices.

With regard to the first criterion (e.g., uniform distribution fitness), the vertices $V(G_2)$ will be compared to an equal number $n(G_3) = n(G_2)$ of vertices $V(G_3)$ that are known to be uniformly distributed along a third and final shell $S(G_3)$ of radius $r(G_3) = r(G_2)$.

The generation of $n(G_3)$ uniformly distributed vertices along a 1D shell (e.g., a circle) is algorithmically simple: divide the circle's 2π radians into $n(G_3)$ equal portions and then use the polar coordinate equations to generate the $n(G_3)$ corresponding vertex positions. The generation of $n(G_3)$ uniformly distributed vertices along a 2D shell (e.g., a thin spherical shell) is not algorithmically simple: an iterative vertex repulsion code [3] was used here to generate $n(G_3)$ roughly uniformly distributed vertices.

The uniform distribution fitness test used here compares the $n(G_2)$ pairs of vertices $V(G_2)_i, V(G_3)_i$ by analyzing the lengths of their corresponding internal non-directed edges $I(V(G_2)_i)_j, I(V(G_3)_i)_j$ (e.g., where $i = \{1, 2, ..., n(G_2)\}, j = \{1, 2, ..., n(G_2) - 1\}$). Some kind of order must be established so that a reasonable correlation exists between $I(V(G_2)_i)_j, I(V(G_3)_i)_j$, and so the lengths of the internal non-directed edges corresponding to each pair of vertices are placed into a pair of sorted bins before the comparison begins

$$L(I(V(G_2)_i)) = \operatorname{sort}[\operatorname{length}[I(V(G_2)_i)_1], \dots, \operatorname{length}[I(V(G_2)_i)_{(n(G_2)-1)}]], \quad (1)$$

$$L(I(V(G_3)_i)) = \text{sort}[\text{length}[I(V(G_3)_i)_1], \dots, \text{length}[I(V(G_3)_i)_{(n(G_3)-1)}]].$$
(2)

Ideally, since $V(G_3)$ are known to be uniformly distributed along $S(G_3)$, the $n(G_3)$ sorted bins $L(I(V(G_3)))$ should all contain identical length distributions (e.g., thus defining a reference distribution $L(I(V(G_3)))_{ref})$. Likewise, if $V(G_2)$ are also uniformly distributed along $S(G_2)$, then the $n(G_2)$ sorted bins $L(I(V(G_2)))$ should also all contain length distributions that are identical to $L(I(V(G_3)))_{ref}$.

The uniform distribution fitness test used here is

$$F_D(G_1) = [0, 1], \tag{3}$$

$$F_D(G_1) = \frac{\sum_{i=1}^{n(G_2)} \sum_{j=1}^{(n(G_2)-1)} \frac{\min[L(I(V(G_2)_i))_j, L(I(V(G_3)_i))_j]}{\max[L(I(V(G_2)_i))_j, L(I(V(G_3)_i))_j]}}{n(G_2)^2 - n(G_2)}.$$
 (4)

It is useful to note that each of a graph's unique internal non-directed edges are analyzed exactly twice throughout the entire test, which is why equation (4) is normalized using $n(G_2)^2 - n(G_2)$, not $(n(G_2)^2 - n(G_2))/2$. With regard to the second criterion (e.g., normal fitness), each external directed edge $E(V(G_1)_i)_j$ corresponds to one intersection vertex $V(G_2)_k$ (e.g., where $i = \{1, 2, ..., n(G_1)\}, j = \{1, 2, ..., n(G_1) - 1\}, k = \{1, 2, ..., n(G_2)\}$. Where both $S(G_1)$ and $S(G_2)$ are centred at the coordinate system origin, the normal fitness test used here is

$$F_N(G_1) = [0, 1], (5)$$

$$F_N(G_1) = \frac{\sum_{i=1}^{n(G_1)} \sum_{j=1}^{(n(G_1)-1)} \frac{\widehat{E}(V(G_1)_i)_j \cdot \widehat{V}(G_2)_k + 1}{2}}{n(G_2)}.$$
 (6)

3 Results

The 1D and 2D shell fitness test results for various $n(G_1)$, $r(G_1)$, and $r(G_2)$ are listed in the following tables

Uniform distribution fitness $F_D(G_1)$ for a 1D shell of radius $r(G_1) = n(G_1)$								
$n(G_1)$	2	4	8	16	32	64	128	
10^{3}	1	0.829572	0.893257	0.946621	0.993124	0.997482	0.998891	
10^{10}	1	0.827916	0.886982	0.93004	0.95885	0.976516	0.986852	
10^{17}	1	0.827916	0.886982	0.93004	0.95885	0.976516	0.986852	

Normal fitness $F_N(G_1)$ for a 1D shell of radius $r(G_1) = n(G_1)$								
$n(G_1)$	2	4	8	16	32	64	128	
10^{3}	1	0.999999	0.999993	0.99997	0.999876	0.999496	0.997962	
10^{10}	1	1	1	1	1	1	1	
10^{17}	1	1	1	1	1	1	1	

Uniform distribution fitness $F_D(G_1)$ for a 2D shell of radius $r(G_1) = n(G_1)$									
$n(G_1)$	2	4	8	16	32	64	128		
10^3	1	0.937087	0.931829	0.974859	0.97905	0.995469	0.998372		
10^{10}	1	0.937088	0.930686	0.973607	0.974738	0.994824	0.997366		
10^{17}	1	0.937088	0.930686	0.973607	0.974738	0.994824	0.997366		

Normal fitness $F_N(G_1)$ for a 2D shell of radius $r(G_1) = n(G_1)$								
$n(G_1)$	2	4	8	16	32	64	128	
10^{3}	1	0.999999	0.999993	0.99997	0.999876	0.999496	0.997963	
10^{10}	1	1	1	1	1	1	1	
10^{17}	1	1	1	1	1	1	1	

4 Discussion

As the fitness test results show, the external directed edges of a complete graph can very nearly form a radially symmetric field at long distance if the number of vertices is great enough. For instance, a 2D shell in 3D space can very nearly reproduce the inverse square law (e.g., field strength proportional to 1/r).

See [2] for the full code and expanded table data. In the full code, the iterative vertex repulsion code [3] has been modified to use the Mersenne Twister pseudorandom number generator [4] in conjunction with the sphere point picking algorithm discussed in [5]. The full code also uses a modified version of the ray-shell intersection code given in [6].

References

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Figure 1: A complete graph G_1 , where $n(G_1) = 5$ vertices $V(G_1)$ (e.g., black disks) are uniformly distributed along a 1D shell $S(G_1)$ (e.g., a gray circle). There are $(n(G_1)^2 - n(G_1))/2 = 10$ internal non-directed edges $I(G_1)$ (e.g., black line segments), and $(n(G_1)^2 - n(G_1)) = 20$ external directed edges $E(G_1)$ (e.g., outward pointing black rays). Where $i = \{1, 2, ..., n(G_1)\},$ $j = \{1, 2, ..., n(G_1) - 1\}$, each vertex $V(G_1)_i$ corresponds to $n(G_1) - 1 = 4$ internal non-directed edges $I(V(G_1)_i)_j$ and $n(G_1) - 1 = 4$ external directed edges $E(V(G_1)_i)_j$.